





























































Why λ calculus works

The process of combining two representations was perfectly uniform. We simply said which of the representations is the functor and which the argument, whereupon combination could be carried out by applying functor to argument and β -converting. We didn't have to make any complicated considerations here.

The load of semantic analysis was carried by the lexicon: We used the λ -calculus to make missing information stipulations when we gave the meanings of the *words* in our sentences. For this task, we had to think accurately. But we could make our stipulations declaratively, without hacking them into the combination process.

What you must do

We have to locate gaps to be abstracted over in the partial formula for our lexical item. In other words, we have to decide *where to put* the λ -bound variables inside our abstraction. For example when giving the representation $\lambda P.P@MARY$ for the proper name 'Mary' we decided to stipulate a missing functor. Thus we applied a λ -abstracted variable to MARY.

We have to decide *how to arrange* the λ -prefixes. This is how we control in which order the arguments have to be supplied so that they end up in the right places after β -reduction when our abstraction is applied. For example we chose the order $\lambda P \lambda Q$ when we gave the representation $\lambda P \lambda Q . \exists x (P @x \land Q @x)$ for the indefinite determiner 'a'. This means that we will first have to supply it with the argument for the restriction of the determiner, and then with the one for the scope.

6.863J/9.611J SP04 Lecture 17



Before we can put λ -calculus to use in an implementation, we still have to deal with one rather technical point: Sometimes we have to pay a little bit of attention which variable names we use. Suppose that the expression \mathcal{F} in $\lambda V \mathcal{F}$ is a complex expression containing many λ operators. Now, it could happen that when we apply a functor $\lambda V \mathcal{F}$ to an argument \mathcal{A} , some occurrence of a variable that is is free in \mathcal{A} becomes bound when we substitute it into \mathcal{F} .

Example

For example when we construct the semantic representation for the verb phrase 'loves a woman', syntactic analysis of the phrase could lead to the representation:

 $\lambda P.\lambda y.(P@\lambda x.LOVE(y,x))@(\lambda Q.\lambda R.(\exists y(Q@(y) \land R@y))@\lambda w.WOMAN(w))$

β-reducing three times yields:

 $\lambda y.(\lambda R.(\exists y(WOMAN(y) \land R@y))@\lambda x.LOVE(y,x))$

Notice that the variable *y* occurs λ -bound as well as existentially bound in this expression. In LOVE(*y*,*x*) it is bound by λy , while in WOMAN(*y*) and *R* it is bound by $\exists y$.





























 $(\lambda P \lambda Q \forall x.(P(x) \rightarrow Q(x)) @ \lambda y.MAN(y)) @ (\lambda R \lambda x R @ \lambda y.LOVE(x, y) @ \lambda P.P(v_1))$

Beta reduced to this...

 $\forall x (\text{MAN}(x) \to \text{LOVE}(x, v_1))$

Now we want to put the "a woman" in front









 $\exists z. \text{THERAPIST}(z) \land \exists y. \text{SIAMESECAT}(y) \land \forall x. ((\text{OWNER}(x) \land \text{OF}(y, x)) \rightarrow \text{LOVE}(x, z)) \\ (A@\text{therapist}) @\lambda z. [(A@\text{s_cat}) \lambda y. [(Every @\lambda x. [OWNER(x) \land \text{OF}(y, x)]) @\lambda x. \text{LOVE}(x, z)]]$

 $*\forall x.OWNER(x) \land OF(y,x) \land \exists y.SIAMESECAT(y) \rightarrow \exists z.THERAPIST(z) \land LOVE(x,z)$

# readings	
number of quantifers	readings
4	14
5	42
6	132
7	429
8	1430
6.863J/9.611J SP04 Lecture 17	















(But the Accord was redesigned for the 2003 model year.)

The roomier, faster, and sleeker sedan's sales stabilized last year,falling by just 1,230 units -- a strong showing in a market that saw combined total passenger car sales fall by 471,000 units.

What's meaning? What's semantics – 2 ends of the spectrum

 Answer 1: whatever it is, it's mapping (translation) between representations
And it depends on all of the tout

And it depends on *all* of the text

- Answer 2: whatever it is, our answer depends on a much more focused task-specific question, viz., <u>information extraction from texts</u>
- Perhaps call this 'natural language <u>engineering'</u>
- These two ends of the spectrum have different characteristics, and difft uses
- · Deep vs. Shallow?

6.863J/9.611J SP04 Lecture 17











































































































Negation
The car doesn't work. ~work(c) vs. not(e) & work'(e,j)
The car almost doesn't work. ALMOST(~work(c)) vs. almost(e1) & not'(e1,e) & work'(e,c) An operator, not first-order
In fact, in not(e), e is really a representative element of a set
John didn't go to class. (yesterday)
To properly interpret negation, we have to figure out that set.
6.8633/9.6113 SP04 Lecture 17


































Intensions vs. Extensions

Last time talking about models, we said the semantics of a unary predicate is the set of all entities in the domain for which the predicate is true

the predicate dog ==> the set of all dogs (in the real world)

This is the **extension** of the predicate dog.

That leaves out possible dogs, future dogs, etc.; makes dog-ness depend only on possibly accidental facts about the real world.

Intension: Map the predicate dog into a mapping from all possible worlds to the set of dogs in that possible world.

the prediate dog ==> [F: possible world w ==> the set of dogs in w

Given a predicate and a possible world, the intension will tell you the set of things that satisfy that predicate in that world.

Intension does a better job of capturing the essence of the concept.

6.863J/9.611J SP04 Lecture 17

