# 6.863J Natural Language Processing Lecture 3: The end of the word 

Instructor: Robert C. Berwick berwick@ai.mit.edu

## The Menu Bar

Administrivia
Lecture 2 posted; Lab 1b("component III") out today

- Finite-state transducer model: the wrapup
- Kimmo \& Laboratory 1b: how-to
- Complexity of fst's - too weak? Too strong? What makes a good computational linguistics representation? A good algorithm?


## The Three Ideas of Two-Level

## Morphology

- Rules are symbol-to-symbol constraints that are applied in parallel, not sequentially like rewrite rules, via sets of transducers
- The constraints can refer to the lexical context, to the surface context or to both contexts at the same time.
- Lexical lookup and morphological analysis are performed in tandem.


## Word of the day

- Turkish word: uygarlas,tiramadiklarimizdanmis,sinizcasina =
uygar+las,+tir+ama+dik+lar+imiz+dan+mis,+siniz+ casina
(behaving) as if you are among those whom we could not cause to become civilized


## Two fs machines in tandem


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## Declarative filtering

- ALL possible string pairs come in
- Spelling change fst's only allow valid spelling pairs through the door (doesn't care about what a valid morpheme is): g:g o:oXs:s
- Morpheme tree allows only valid root + affix sequences through the door (doesn't care about spelling): beer be+er
- Combination admits only permissible lexical:surface word forms





## Ambiguities

"upper language"



## Insert 'e' before non-initial $z, \mathrm{~s}, \mathrm{x}$ ("epenthesis")



## Successful pairing of foxes,fox+s



| f | $\bigcirc$ | X | 0 | e | S | \# |  | rface |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| F | 0 | X |  | d | S | \# |  | exical |

## Can a transducer do this?

- igpay atinlay is unfay
- define Cons [ b ] ; define Vowel [a|e|i|o|u]; define Ltr [ Cons | Vowel ]+ ; define Limit [" " | "\t" | .\#. ];
- Suppose just one consonant, e.g, be
- Strategy?




## Constraints on both sides



N:m correspondence requires a following $p$ on the lexical side.
In this context, all other possible realization of a lexical N are prohibited.

p:m correspondence requires a preceding m on the surface side.
In this context, all other possible realization of a lexical $p$ are prohibited.

## Parallel application - how?


kamanan man
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## Sequential Application

kan p a n


## Machine Rule 1 ("N goes to m")

## Rule 1: $\mathbf{N} \rightarrow \mathbf{m} \mid \ldots \mathbf{p}$


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## Machine Rule 2 ("p goes to m")

Rule 2: $\mathrm{p} \rightarrow \mathrm{m} \mid \mathrm{m}$ $\qquad$

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## Sequential Application in Detail

kanpan

| 0 | 0 | 0 |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 2 | 0 | 0 | 0 |

kampan
$\begin{array}{llll}0 & 0 & 0 & 1 \\ & 0 & 0 & 0\end{array}$
kamman


## OK then...

- To apply all 5 rules at once, we can simply intersect them all, right? (So given character pair has to pass all the spelling checks at each point)
- Right?
- Wrong!


## Relations

Ordered Set: members are ordered.
Ordered Pair: <A,B> vs. <B,A>
Relation: set whose members are ordered pairs - (lexical, surface)

## Relations

An Infinite Relation:.


Identity Relation: <"fly", "fly">

## What does transducer define?

- Lexical form is a finite-state language (a regular language)
- Surface form is a regular language
- Transducer pairs (lexical form, surface form) $=$ regular relation ("rational relation:")
- What are the properties of rational relations? Is this the same as what Kimmo does?


## Closure properties of FSTs not same

 as FSAs!- NOT closed under intersection
why? Example: $(a: b)^{*}(0: c)^{*} \cap(0: b)^{*}(a: c)^{*}$
Intuition: $1^{\text {st }}$ makes equal $a^{\prime} s, b$ 's; $2^{\text {nd }}$, equal $a^{\prime} s$ \& C's
So output is: < , >, <a, bc>, <aa, bbcc>, ...
- NOT possible to make FTNs deterministic in general either...
why? Consider following FST:

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## Composition

Composition is an operation on two relations.
Composition of the two relations $\langle x, y\rangle$ and <y,z> yields <x, z>

Example: <"cat", "chat"> with <"chat", "Katze"> gives <"cat", "Katze">

## Composition

- A c B The relation $C$ such that if $A$ maps $x$ to y and B maps y to $\mathrm{z}, \mathrm{C}$ maps x to z .


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Merging the two networks

## Composition



The Composition of the Networks

## What is this reminiscent of?

## Composition

kanpan
$0 \quad 0 \quad 0|3| \begin{array}{lll}0 & 0 & 0\end{array}$
kamman


FTNs ARE closed under composition

## What are the implications?

- FTNs inherently require backup if simulated (in the worst case) - Kimmo at least NP-hard (proof later on)
- Empty elements cause computational complexity (unless restricted - equal length condition)
- Composition can save us, but then rule ordering must be watched carefully


## Sequential Application in Detail

$$
\begin{array}{ccccccc}
k & a & N & p & a & n \\
0 & 0 & 0 & 2 & 0 & 0 & 0 \\
k & a & m & p & a & n \\
& 0 & 0 & 0 & 1 & 0 & 0
\end{array} 0
$$

```
            Let's see an example - English
            races'
_Recognizing surface form "races'".
    0 (r.r) --> (1 1 1 1 2 2 1 1)
                EP G Y EL I
    1 (a.a) --> (1 1 1 4 4 1 2 1)
                                EP G Y EL I
    2 (c.c) --> (1 2 16 2 11 1)
    3
        (e.0) --> (\begin{array}{llllllll}{1}&{1}&{16}&{1}&{12}&{1}\end{array})
                                EP G Y EL I
4
                Entry |race| ends --> new lexicon N, config( (\begin{array}{lllllll}{1}&{1}&{16}&{1}&{12}&{1}\end{array}) EP G Y EL I
- note that each of the 5 automata are running
"in parallel" (what is the firs3 first one?)
```

```
Backup search
Problem; e is paired with 0 (null)...!
    (which is wrong - it's guessing that the form is
    "racing" - has stuck in an empty (zero) character
    after c but before e) - elision automaton has
    2 choices
    This is nondeterminism in action (or inaction)!
    5
    Entry /O ends --> new lexicon C1, config((\begin{array}{llllll}{1}&{1}&{16}&{1}&{12}&{1}\end{array})
                                    EP G Y EL I
6
5
6
5
4
    Entry /0 is word-final --> path rejected (leftover input).
        (+.0) --> ((1)1 1 16 1. 1 13 1)
        EP G Y EL I
    Nothing to do.
    (+.e) --> automaton Epenthesis blocks from state 1.
    Entry |race| ends --> new lexicon P3, config((\begin{array}{lllllll}{1}&{1}&{16}&{1}&{12}&{1}\end{array})
                                    EP G Y EL I
```


## 22 steps later

```
(e.e) --> (1 1161414 1)
EP G Y EL I
    Entry |racel ends --> new lexicon N, ( \(\left.\begin{array}{llllll}1 & 1 & 16 & 1 & 14 & 1\end{array}\right)\)
                                    E G Y EL I
    Entry /0 ends --> new lexicon C1, config ( \(\left.\begin{array}{llllll}1 & 1 & 16 & 1 & 14 & 1\end{array}\right)\)
    Entry /0 is word-final -->rejected (leftover input)
        (+.0) --> ( \(\left.\begin{array}{llllll}1 & 1 & 16 & 1 & 15 & 1\end{array}\right)\)
        (s.s) --> ( \(\left.\begin{array}{llllll}1 & 4 & 16 & 2 & 1 & 1\end{array}\right)\)
            Entry +/s ends--> new lexicon C2, ( \(\left.\begin{array}{llllll}1 & 4 & 16 & 2 & 1 & 1\end{array}\right)\)
            Entry /O is word-final -->rejected(leftover input)
            ('.') --> (1 116111 )
                End --> lexical form ("race+s'" (N PL GEN))
```

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## Laboratory 1b

- Goals:

How to use_Kimmo to analyze_another language (Spanish),_as example "front end"

- Build automata for some simple Spanish morphological/phonemic rules (that interact)
- Build lexicon
- Learn what is hard and what is easy about this
- Recognize all and only the words in spanish.rec; Generate all the surface forms
- Resources:
- Lab1b pdf file link from web page
- File of all the surface words to parse/reject (covering the phenomena) spanish.rec, also linked from web page
- Pykimmo, pcimmo \& documentation
- Program to `compile' rules into automata: fst


## What you must turn in (via URL)

A description of how your system operates
2. URL ptrs to your .lex and .rul files
span.lex
span.rul
3. A $\log$ of a recognition run on the file spanish.rec which is linked on the web page \& also at toplevel on course locker
4. Discussion of what you built/why
5. You must answer 3 questions:

## The questions <br> - What is your name? <br> - What is your quest? <br> . What...

## The phenomena under study

You are given the ofthography, inctuding some speciat characters to stand for the accented ones á,é,ó,ü,ñ ; and some underlying characters you may find essential, such as J, C, Z.

- Wise to proceed by first building the automata (rul) file; then the lexicon(s) - because you can test the rules without any lexicon by generation of a surface form
- The automata can be built (roughly) by considering each phenomenon separately
- 3 kinds of phenomena


## The phenomena

1. $g$-j mutation
2. z-c mutation
3. pluralization
4. Noun endings
5. Verb conjugation - 1 form

## Phenomenon 1: g-j mutation

g-j mutation
$g \rightarrow j$ before a back vowel
coger (catch, infinitive); cojo, coges, coge, cogemos, cogen, coja NOTE: coger is NOT the lexical underlying form!!!)

- But some verbs not subject to this (exceptions!) llegar (arrive); llego, llegan, pagar (pay); pago, pagan
- Don't acccept */lejo, *lleja, *cogo, *coga (the words don't come marked with * on their sleeves, of course!)
- Hint: can use the lexical (underlying) character J to solve (but there are other ways to do it)


## How to build Kimmo systems

How to build lexicons using morpheme states and the actual lexical entries How to build automata for spelling changes

## Format for .lex file - 2 parts

1. Lexicon: Morpheme classes - name all the states, some transitions
(1 or more blank lines)
2. Lexicon entries: transitions between the states
(Recall that we consider only the underyling form combinations here - stems + affixes, not spelling change rules on the surface)

## Example: lexicon design

Phenomena: Nouns and Verbs take different endings

Answer:
Different morpheme states for Nouns and Verbs

## Example surface (s) underlying (u) pairs tell us what to do

- ciudad
ciudad [N(city)]
- ciudades ciudad+s [N(city)pl]


Q: what do we need to add to noun alternation?


## How do we build spelling change automata?

## Example: look at phenomenon, then

 see first how to describe- What is the left and right context of the change?
- Write it as a declarative constraint
- Remember that you can use both the surface and the lexical characters to admit or to rule out a possibility


## Phenomenon 2: z-c mutation

- z-c mutation
$z \rightarrow c$ before front vowels, $z$ otherwise cruzar (to cross); cruzo, cruzas, cruza, cruzamos, cruzan, cruce
- If $s$ causes a front vowel (e.g., e) to surface, then the rule still applies: lápiz, lápices (pencil, pencils) [ /^piz, /^pices]

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## Now add the arcs...

## Phenomenon 3: pluralization

- Adding $s$ to a noun that ends in a consonant forces a surface $e$ to appear: ciudad (city); ciudades
- This can interact with other rules, e.g., z-cmutation:
lápiz, lápices
- Nouns ending in a vowel are not subject to this rule: bota, botas


## The lexicon - take 2

Add a gloss at the very end of the process, so as to return the feature list and 'translation', e.g., venzo [1p sg pres indic conquer, defeat] (first person, singular, present tense, indicative)

- We'll show how to add this in a moment
- You will deal with two types of 'endings'

1. Noun endings: plural suffix $+s$
2. Verb endings: verb stem + tense markers

Simplest: infinitive marker +ar, +er, +ir
See table in pdf file for details: $5 \times 3$ table for Present tense; ditto for Subjunctive tense ("I might....)

## Instead of writing fst tables...

- You can use the program fst
- To run:
build fst type rules in file spanish.fst, then
fst -o ~yourpath/spanish.rul ~yourpath/spanish.fst
- Also script to print fst files to dot, for ps viewing
- Format for fst rules:


## FST rules

" $b$ after a vowel turns to $a$ " subset vowel a e
machine "bintoa"
state foo
vowel:vowel bar
b:b foo
c:c foo
d:d foo
others reject
rejecting state bar
b:e foo
b:b reject
others foo

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## Design of morpheme machine

- One big fsa
- Like this...


- List states - in cyan - followed by sets of transition labels (possible outgoing arcs)
Begin:
N_root:
$\begin{array}{ll}\text { N_Root1 } & \text { Adj_Roo } \\ \text { Poss } & \text { To_adj }\end{array}$
V_Root
------------------------------------------------
- For each arc: List transitions \& next states, and output N_root1:
Kol N_root
Kitab N_root

Noun['arm']
Noun['book']
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## The End

End:
0 \# ""
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