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LINEAR RULES

The simultaneous rule model is, of course, just one of a number of conceivable formalisms that would impose a finite-state restriction on phonological rules. One could, for example, propose that phonological rules be finite transducers in the literal sense. No one would take such a suggestion seriously because of the linguistic inappropriateness of the formulations it would require. We will, however, give serious consideration to another type of rule, so far employed very rarely in generative phonology, which shares some of the properties of the iterative type but which, like the simultaneous rule, is a finite-state device in its mapping power. This new type of rule, to be called linear, has two subvarieties which we will refer to as the right-linear and the left-linear.

We consider first right-linear rules. To describe how these function it will be convenient to use the notion of right-overlap. If $M = P \rightarrow Q/R - S$ and $N = P' \rightarrow Q'/R' - S'$ are elementary rules, we will say that *M* right-overlaps *N* if and only if

- (i) $RPS = R'P'S'$,
- (ii) *R* is not longer than *R'*, and
- (iii) *RP* is longer than *R'*.

Thus *M* right-overlaps *N* if

$M = ushataraa \rightarrow ushataraa/ - naam,$
 $N = n \rightarrow \eta/ us - ataranaanam$

or if

$M = nata \rightarrow nata/ - pikappikai$
 $N = a \rightarrow \acute{a}/ nat - pikappikai$

However, *M* does not right-overlap *N* if

$M = us \rightarrow us/ - nataraanaan,$
 $N = n \rightarrow \eta/ us - ataranaan$

or if

$M = watumis \rightarrow watumis/ - ?as,$
 $N = \emptyset \rightarrow a/ watumis - ?as$

We now formalize right-linear rules as in (41).

(41) (a) A right-linear rule is an expression of the form $R:X$

where *X* is a schema subsuming elementary rules
 subrules of $R:X$ and *R* is a constant denoting the
 mode of application next to be defined. There is no
 restriction on the form of subrules.

(b) Let *P* and *Q* be phonological strings and let $R:X$ be a
 right-linear rule. Then $R:X$ maps *P* into *Q* if and only
 if there is a sequence $((P_1, Q_1), \dots, (P_n, Q_n))$ of ordered
 pairs of strings such that

- (i) *n* is odd;
- (ii) $P = P_1 \dots P_n$;
- (iii) for each even *i*, $2 \leq i \leq n-1$, the elementary
 rule

$P_i \rightarrow Q_i/Q_{i-1} - P_{i+1} \dots P_n$
 is a subrule of $R:X$;

(iv) for each odd *i*, $1 \leq i \leq n$, $P_i = Q_i$ and the
 elementary rule

$P_i \rightarrow Q_i/Q_{i-1} - P_{i+1} \dots P_n$
 does not right-overlap any subrule of $R:X$;

(v) $Q = Q_1 \dots Q_n$.
 The sequence $((P_1, Q_1), \dots, (P_n, Q_n))$, which we will
 write also as

$P_1 \quad P_2 \quad P_3 \dots P_{n-1} \quad P_n,$
 $Q_2 \quad Q_3 \dots Q_{n-1}$

will be referred to as an application of $R:X$ with input *P*
 and output *Q*. The *i*th rightward step of the application
 is defined as the elementary rule $P_i \rightarrow Q_i/Q_{i-1} \dots Q_{i-1}$

— $P_{i+1} \dots P_n$. Thus the even-numbered steps are subrules of $R:X$ and the odd-numbered steps right-overlap no subrules of $R:X$.

The Sanskrit nasal retroflexion rule can be used again to illustrate the linear modes of application. Using the same schema as we did in the simultaneous rule (30), we can construct the right-linear rule (42).

$$(42) \quad R: \begin{bmatrix} +nas \\ +cor \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} -ant \\ (\sigma) \end{bmatrix} / \$^* \begin{bmatrix} +cor \\ +cont \\ -ant \\ -distr \end{bmatrix} [-cor]^* - [+son] \*$

This will give the same results as the simultaneous rule. In particular, it still convert *usnataranaam* into *usnataranaam* by virtue of the same application, namely:

$$u\$ \quad \begin{matrix} n \\ \eta \end{matrix} \quad \begin{matrix} atarana \\ \eta \end{matrix} \quad \begin{matrix} n \\ \eta \end{matrix} \quad aam$$

Here, however, the application is regarded as consisting of the following rightward steps:

1. $u\$ \rightarrow u\$/-nataranaam$
2. $n \rightarrow \eta/u\$/-ataranaam$
3. $atarana \rightarrow atarana/u\$/-naam$
4. $n \rightarrow \eta/u\$/natarana-aam$
5. $aam \rightarrow aam/u\$/natarana-$

Notice that the dash appears in the leftmost possible position in the first step and occupies positions successively farther to the right in subsequent steps. Furthermore the output of each step other than the last is the input to the next step. Consider in particular the even-numbered steps, which are also subrules of (42). Since the other steps are vacuous, the input to the application is also the input to the second step, the output of the second step is the input to the fourth step, and the output of the fourth step is the output of the application as a whole. Now observe that 2

and 4 are precisely those subrules which were invoked in application (27) of the iterative version of the rule (26). Thus the right linear rule, though formally very similar to the simultaneous rule, is in another respect quite like the iterative rule. In general, the function of R is to restrict iterative applications to those which proceed through the input string in a strictly left-to-right manner, stopping when the right end of the string is reached even if the string in its current form is the input to a further subrule. These conditions on application are sufficient to reduce the power of iterative rules to that of finite-state devices. For it is possible to demonstrate (43).

- (43) If N is a monogenic right-linear rule there is a left-transducer M_1 and a right transducer M_2 such that (M_1, M_2) provides each nonempty phonological string with exactly the same realization that N does.

* We can let $N = R:X$. As in the proof of (35) we can assume that $X = \{X_1, \dots, X_m\}$ is a primitive schema in which each X_i has the form $Y_1 \rightarrow S_i/U_1 - V_i$, where S_i is a phonological string and Y_1, U_1 , and V_1 subsume only phonological strings.

Suppose first that N has no subrules of the form $\emptyset \rightarrow P/Q - R$. The construction of (M_1, M_2) proceeds exactly as the proof of (35) down through the construction of M_1 . M_2 has the states, input alphabet, and output alphabet described in the proof of (35), and a similar designation of left and right terminal states, but a different set of instructions. Specifically, M_2 will have, for each state (u, y, k) and each input symbol (w, v, Δ) , the instruction $(u, y, k)(w, v, \Delta) \rightarrow Q(u', y', k')$ where all the conditions of (44) hold.

- (44) (a) For each r , $1 \leq r \leq m$, $(u')_r$ is the rightmost state of that computation of U_r which has $(u)_r$ as its leftmost state and Q as its input. (Notice that u' must be identical to u if $Q = \emptyset$.)
 (b₁) (Identical to (36b₁))
 (b₂) (Identical to (36b₂))

The proof that (M_1, M_2) does the same work as N will not be given, since it follows the same general lines as the corresponding portion of the proof of (35). The minor differences that exist between the two proofs stem from the discrepancy between (44a) and (36a).

Suppose now that N has some subrules of the form $\emptyset \rightarrow P/Q-R$.

Consider a modified formalism that allows the symbol Z , distinct from all phonological units and not subsumed under S , to be referred to in rules. Obtain Y'_1, U'_1 , and V'_1 from Y_1, U_1, V_1 , respectively, by substituting Z^*AZ^* for each phonological unit A . Let $X' = \{X'_1, \dots, X'_m\}$, where

$$X'_1 = Y'_1 \rightarrow S_i/U'_1 - V'_1, \text{ if } Y_1 \text{ does not subsume } \emptyset,$$

$$X'_1 = \left\{ \begin{array}{l} \left[Y'_1 \rightarrow S_i/U'_1 - V'_1 \right] \\ \left[Z, S \right]^* S \rightarrow S^* \{ Z, S \}^* - \{ Z, S \}^* \end{array} \right\} \text{ if } Y_1 \text{ sub-} \\ \text{sumes } \emptyset, \\ Z \rightarrow S_i/U'_1 - V'_1$$

Let X'' be a primitive schema equivalent to X' . The rule $R: X''$ has no subrules of the form $\emptyset \rightarrow P/Q - R$ but will simulate N in the following sense. Where N maps a nonempty phonological string $A_1 \dots A_n$ into $B_1 \dots B_n$, $R: X''$ will map $ZA_1Z \dots A_nZ$ into $K_1B_1 \dots K_nB_nK_{n+1}$, where each K_i is a string (possibly empty) of Z 's. In the manner described in the preceding paragraph we can construct a left transducer M_1 and a right transducer M_2 such that (M_1, M_2) performs the same mapping as N' . It is a trivial matter to construct a left transducer M_0 that inserts Z between each pair of adjacent segments in a nonempty input string and also at the beginning and end of the string. Even more trivial is the construction of a transducer M_3 that deletes all occurrences of Z . It is obvious that the 4-tuple (M_0, M_1, M_2, M_3) will provide each nonempty phonological string with the same realization that N does. By virtue of some results of Schützenberger (1961), this 4-tuple is equivalent to a single finite transducer and hence to an ordered pair of machines the first of which is a left transducer and the second a right transducer.*)

In formalizing simultaneous rules we failed to accommodate insertion processes as they are usually formulated. Insertion

subrules of right-linear rules are interpreted in the desired way, however. The rule of vowel-doubling proposed for Mohawk by Postal (1969) can, for example, be given as the right-linear rule (45).

$$(45) \quad R: \emptyset \rightarrow \sigma / S^* [-sy] \rightarrow \begin{bmatrix} +sy \\ \sigma \end{bmatrix} S^*$$

This will turn *watunis²as* into *watunisa²as* by virtue of the application

$$\text{watunis } a^2 \text{as}$$

This application has the three steps

$$\begin{array}{l} \text{watunis} \rightarrow \text{watunis}' / -^2 \text{as} \\ \emptyset \rightarrow a / \text{watunis} -^2 \text{as} \\ ^2 \text{as} \rightarrow ^2 \text{as} / \text{watunisa} - \end{array}$$

Left-linear application is the mirror image of right-linear application. Corresponding to the notion of right-overlap is that of left-overlap. If $M = P \rightarrow Q/R - S$ and $N = P' \rightarrow Q' R' - S'$ are two elementary rules we will say that M left overlaps N if and only if

- (i) $RPS = R'P'S'$,
- (ii) S is not longer than S' , and
- (iii) PS is longer than S' .

Then we formalize left-linear rules as in (46).

- (46) (a) A left-linear rule is an expression of the form $L: X$ where X is a schema subsuming elementary rules (the subrules of $L: X$) and L is a constant denoting the mode of application next to be defined.

(b) Let P and Q be phonological strings and let $L: X$ be a left-linear rule. Then $L: X$ maps P into Q if and only if there is a sequence $((P_1, Q_1), \dots, (P_n, Q_n))$ of ordered pairs of strings such that

- (i) n is odd;
- (ii) $P = P_n \dots P_1$;

(iii) for each even i , $2 \leq i \leq n-1$, the elementary rule

$$P_i \rightarrow Q_i/P_n \dots P_{i+1} - Q_{i-1} \dots Q_1$$

is a subrule of $L:X$;

(iv) for each odd i , $1 \leq i \leq n$, $P_i = Q_i$ and the elementary rule

$$P_i \rightarrow Q_i/P_n \dots P_{i+1} - Q_{i-1} \dots Q_1$$

does not left-overlap any subrule of $L:X$;

$$(v) Q = Q_n \dots Q_1.$$

The sequence $((P_n, Q_n), \dots, (P_1, Q_1))$, which we can write as

$$P_{n-1} \quad P_{n-2} \quad P_2 \quad P_1 \\ Q_{n-1} \quad Q_2 \quad Q_2 \quad P_1$$

will be called an application of $L:X$ with input P and output Q . By the i th leftward step of this application we mean the elementary rule

$$P_i \rightarrow Q_i/P_n \dots P_{i+1} - Q_{i-1} \dots Q_1$$

The Sanskrit nasal retroflexion can be expressed by a left-linear rule involving the same schema as was used in (42). The left-linear rule gives the same results as the right-linear one, and its application has the same form. For example *usnatarāṅānam* is processed by the application

$$us \quad n \quad ataraa \quad n \quad aam \\ \eta \quad \eta \quad \eta \quad \eta$$

Here, however, the application is regarded as consisting of the following leftward steps:

1. aam \rightarrow aam/usnatarāa—
2. n \rightarrow n/usnatarāa—aam
3. ataraa \rightarrow ataraa/usn—ānam
4. n \rightarrow n/us—atarāṅānam
5. us \rightarrow us/—natarāṅānam

This application bears the same relation to the iterative application (28) as the right-linear application of (42) displayed earlier bears to (27).

By a proof symmetric to that outlined for (43) we can show that for each monogenic left-linear rule N there is a right-transducer M_1 and a left-transducer M_2 such that (M_1, M_2) provides each phonological nonempty phonological string with the same realization that N does.

Let us now review some of the formal results concerning simultaneous and linear rules. We have seen that each monogenic rule that is right-linear, left-linear, or simultaneous is equivalent in the mapping it effects to an ordered pair of machines each of which is right transducer or a left transducer. Hence each such rule is equivalent to some finite transducer in the sense of Schützenberger. By a series of simple constructions, here omitted, we can show converses to these assertions. Specifically, every finite transducer can be simulated by some simultaneous rule.¹ Furthermore, every finite transducer can be simulated by each of at least four ordered pairs of linear rules, representing the four possible combinations of directionality:

| | |
|-----------------|-----------------|
| <i>1st rule</i> | <i>2nd rule</i> |
| right-linear | right-linear |
| right-linear | left-linear |
| left-linear | right-linear |
| left-linear | left-linear |

A further result, due to Schützenberger (1961), is that any n -tuple of finite transducers operating in sequence can be reduced to a single finite transducer. Consequently, if f is a many-to-one mapping of phonological strings into phonological strings, the following assertions are entirely equivalent:

- (a) f can be effected by a sequence of finite transducers;
- (b) f can be effected by a finite transducer;

1 It would be absurd, of course, to maintain that the class of phonological rules coincides with the class of rules that can perform finite-state transductions. As the formalism stands it is still possible to formulate many implausible rules, even in a simple way. The extent to which the further necessary restrictions can be stated in general formal terms is not a question we will go into here. It is obvious, however, that substantive constraints, such as those imposed by marking conventions, are indispensable.

- (c) f can be effected by a sequence of simultaneous rules;
 (d) f can be effected by a simultaneous rule;
 (e) f can be effected by a sequence of right-linear rules;
 (f) f can be effected by a sequence of left-linear rules.

It makes no difference to mapping capacity, then, which of the three types of rules (simultaneous, right-linear, left-linear) we allow in phonological descriptions. We will therefore have to judge the three types solely on the basis of the formulations they yield. In general, we prefer formulations which reflect naturalness, plausibility, or significant generality with corresponding notational economy. From the point of view of this criterion, it seems clear that we need a formalism that allows both right-linear and left-linear rules, and we shall consider a variety of cases that support this view.

Southern Paiute has a well-known rule which, counting from left to right, stresses the even-numbered nonfinal vowels of a word. If we follow Chomsky and Halle (1968: 244-9) concerning the nature of underlying representations in Southern Paiute, assuming in particular that underlying vowels are all unstressed and are to be equated with Sapir's moras, and if we assume that their rule (44) (which among other things deletes word-final consonants) precedes the alternating stress rule, then we can formulate the latter rule in right-linear fashion as in (47) below. (This formulation is similar to DRULE R7 of Bobrow and Fraser (1968: 769).)

$$(47) \quad R: \begin{bmatrix} +\text{syll} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} +\text{stress} \\ (\sigma) \end{bmatrix} / \$* \begin{bmatrix} +\text{syll} \\ -\text{stress} \end{bmatrix} [-\text{syll}]^* - \$\$*$$

To see how this rule works consider the word *natpikkappikai* 'they threw rocks at one another'. The underlying form given here is that presumably imposed by the analysis of Chomsky and Halle. This word is the input of one application of (47), namely:

nat^a a^a pikk^a a^a pikk^a a^a iⁱ

This application consists of the following rightward steps:

1. nat → nat/—apikkappikai
2. a → á/nat—pikkappikai
3. pikk → pikk/natá—appikai
4. a → á/natáppikk—ppikai
5. ppik → ppik/natáppikká—ai
6. a → á/natáppikkáppik—i
7. i → i/natáppikkáppiká—

The even-numbered steps are subrules of (47) and the odd-numbered steps are vacuous and do not right-overlap any subrule of (47). Note that although (47) has the subrules

i → i/natap—kkappikai
 i → i/natappikkapp—kai

it will never yield the incorrect stress pattern **natáppikkappikai* by virtue of any putative application

natapⁱ iⁱ kkappⁱ iⁱ kai

The first step here would be

natap → natap/—ikkappikai,

which right overlaps the subrule

a → á/nat—pikkappikai

The output *natáppikkáppikáí* is subject to further rules which convert it ultimately into *naráwíkkáppixáá*, where . denotes voicelessness. (See Chomsky and Halle, loc. cit. and also Harris 1966b. We assume that the 'spirantization' rule (46) of Chomsky and Halle follows rather than precedes the alternating stress rule; apparently the two rules can occur in either order with no difference in effect.)

Notice that we cannot obtain a simultaneous or left-linear version of the Southern Paiute stress rule merely by substituting L or S for R in (46); if we did we would have a rule that stressed

every nonfinal nonlast vowel in a word, deriving forms like **nināpkkōppkkii*. Apparently the optimum solution in simultaneous or left-linear mode is something like formulation (48).

$$(48) \quad L, S: \left[\begin{array}{c} +syl \\ \sigma \end{array} \right] \rightarrow \left[\begin{array}{c} +stress \\ (\sigma) \end{array} \right] / [-syl] * [+syl] \\ [-syl] * [+syl] [-syl] * [+syl] * [-syl] * - \$ *$$

This rule states explicitly that a vowel to be stressed must be preceded by an odd-number of vowels within the word, whereas the right linear rule (47) requires merely that the nearest preceding vowel be unstressed. This difference in the way the left-hand context is described is the source of the somewhat greater simplicity of the right-linear formulation.

The case just considered represents a common phenomenon. Eastern Ojibwa, as described by Bloomfield (1956), apparently has an alternating stress rule much like that of Southern Paiute, though with some additional complications. The distinction between underlying long and short vowels, irrelevant in Southern Paiute where phonetic long vowels seem to derive from underlying geminate clusters, is crucial in Eastern Ojibwa. According to Bloomfield (p. 5) the odd-numbered vowels in a sequence containing only short vowels are reduced in "loudness" sometimes to the point of complete disappearance and undergo changes in quality. Other vowels, whether long or short, are not reduced. Furthermore, the last vowel in a word is never reduced. If we interpret reduced loudness as lack of stress and nonreduced loudness as presence of stress, and if we assume underlying vowels to be unstressed, we can formalize Bloomfield's descriptive statement (apart from the quality changes) in terms of the sequence of rules in (49). In (49a) and (49b) any of the three application modes, right-linear, simultaneous, and left-linear, are possible, and we can omit the mode designator.

$$(49) \quad (a) \quad \left[\begin{array}{c} +syl \\ +uns \\ \sigma \end{array} \right] \rightarrow \left[\begin{array}{c} +stress \\ (\sigma) \end{array} \right] / \$ * - \$ *$$

$$(b) \quad \left[\begin{array}{c} +syl \\ \sigma \end{array} \right] \rightarrow \left[\begin{array}{c} +stress \\ (\sigma) \end{array} \right] / \$ * - [-syl] * \\ (c) \quad R: \left[\begin{array}{c} +syl \\ \sigma \end{array} \right] \rightarrow \left[\begin{array}{c} +stress \\ (\sigma) \end{array} \right] / \$ * \left[\begin{array}{c} +syl \\ -stress \end{array} \right] [-syl] * - \$ *$$

(49a) stresses long vowels, (49b) stresses final vowels, and (49c), the only essentially right-linear rule, stresses the even-numbered vowels in a sequence of syllables containing vowels still unstressed (which are identical to nonlast short vowels after application of (49a-b)). We can illustrate the operation of (49) with part of the paradigm of a verb meaning 'to arrive'.

| | | | |
|-------------|--------------|---------------|----------------------------------|
| ninakoššin | takoššin | takoššino:k | 48a |
| ninakoššin | takoššin | takoššino:k | 48b |
| ninakoššin | takoššin | takoššino:k | 48c |
| nentakuššin | tekoššin | tekoššeno:k | reduced-vowel quality changes |
| 'I arrive' | 'he arrives' | 'they arrive' | |

The best we can do in trying to formulate the Eastern Ojibwa alternating stress rule in simultaneous or left-linear mode is apparently (50).

$$(50) \quad L, S: \left[\begin{array}{c} +syl \\ \sigma \end{array} \right] \rightarrow \left[\begin{array}{c} +stress \\ (\sigma) \end{array} \right] / (\$ * [+stress]) * [-syl] * \\ [+syl, -stress] ([-syl] * \left[\begin{array}{c} +syl \\ -stress \end{array} \right] [-syl] * \\ [+syl, -stress]) * [-syl] * - \$ *$$

(50) says explicitly that a vowel is to be stressed if an odd-number of unstressed vowels intervene between this vowel and the nearest preceding stressed vowel or, if there is no such stressed vowel, the beginning of the word.

Notice that the optimal right-linear versions of the Southern Paiute and Eastern Ojibwa alternating stress rules, (47) and (49c), are almost identical, whereas the optimal simultaneous (and left-linear) versions, (48) and (50), diverge because of the extra

[—stress] specifications necessary in the left-hand context of (50). Thus the right-linear formulations reveal more clearly that the two languages exploit essentially the same alternating stress rule, despite their different treatment of long vowels and final-syllable vowels.

Alternating stress is apparently just one manifestation of a widely attested process which in some sense strengthens or weakens alternate vowels in certain kinds of sequences. We consider briefly three cases in which vowel strength is manifested as length, voicing, and retention, respectively.

In Tübatulabal, according to Swadesh and Voegelin (1939), there is a word-level rule which lengthens the first vowel and each even-numbered vowel in a sequence containing only short vowels.² However, a short vowel will not be lengthened if the next following vowel, if any, is long. Some input-output relations defined by this rule are given below. We transcribe the forms according to the principles followed by Swadesh and Voegelin, inferring from one of their suggestions that \exists (a vowel-shortening variety of glottal stop) is the final segment of the morpheme *la* \exists 'going'.

| | | |
|-----------------------------|------------------------------|---------------------|
| dawaginana $\lambda\exists$ | adawaginana $\lambda\exists$ | (input) |
| da:wagi:nana:la \exists | a:dawa:gina:nala: \exists | (vowel lengthening) |
| ta:wagi:nana:la | a:dawa:gina:nala | (later rules) |
| 'he goes along | 'he went along | |
| causing him | causing him | |
| to see' | to see' | |

In Japanese, according to Han (1962: 36-43), alternate vowels are devoiced in a sequence of the form $C_1A_1\dots C_nA_nC_{n+1}$ where the C_i are voiceless stops and the A_i are short high unaccented vowels. Whether it is the odd or even vowels that are devoiced is a matter of free variation. Thus we have *pukupukuto* and *pukupukuto* as freely varying pronunciations of *pukupukuto*.

According to Delattre (1951: 348), French has a rule which deletes the even-numbered vowels in a sequence of the form

² McCawley (1969) has discussed this rule in great detail and proposed a left-to-right iterative formulation that could be regarded as right-linear.

$C_1A_1\dots C_nA_nC_{n+1}$ where the C_i are consonants, A_i is any vowel, and A_2 through A_n are all schwas. For example, orthographic *elle ne me le redemande pas*, which presumably has the phonological shape *elanemalaradamāpasa* just before the schwa deletion rule is to apply, is pronounced in Delattre's norm as *elanlaradamāpa*.

The three rules just described bear a striking formal resemblance to left-to-right alternating stress rules and seem, like them, to be most naturally expressed by right-linear formulations.

So far we have considered only rules that affect alternate vowels from left to right. Rules that affect alternate vowels from right to left seem to be somewhat rarer, but they do, in fact, exist, and provide evidence that we must allow left-linear application in phonology. We again consider the Tübatulabal language, which (according to Voegelin 1935: 75-78) has the following block of word stress rules:

- (i) the last vowel is stressed,
- (ii) each long vowel is stressed, and
- (iii) in a sequence containing vowels not stressed by (i) or (ii), the even-numbered vowels, counted from right-to-left, are stressed.

This block of rules is later than the left-to-right alternating length rule that was mentioned above, and indeed follows certain vowel shortening rules which obscure the alternating length pattern. To formalize (i)-(iii) we can look to the very similar Eastern Ojibwa stress rules (49). (i) and (ii) can simply be identified with (49a) and (49b). A formal version of (iii) can be obtained from (49c) by simply replacing *R* with *L* and replacing the contextual portion of (49c) by its mirror image. The stress pattern of Tübatulabal is given, then, by (51).

- (51) (a) (like (49a))
- (b) (like (49b))
- (c) $L: \left[\begin{array}{c} +syl \\ \sigma \end{array} \right] \rightarrow \left[\begin{array}{c} +stress \\ (\sigma) \end{array} \right] / \$* \left[\begin{array}{c} -syl \\ -stress \end{array} \right] \$*$

Examples of derivations with (51) are given below.

| | | |
|------------|-------------------|---------|
| witaghatal | witaghatala:batsu | (input) |
| witaghatal | witaghatala:batsi | (51a) |
| witaghatal | witaghatala:batsü | (51b) |
| witaghatal | witaghatala:batsü | (51c) |

The optimal simultaneous (or right-linear) version of the Tübatulabal alternating stress rule seems to require the much more complex schema (52).

$$(52) \begin{bmatrix} +syl \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} +stress \\ (\sigma) \end{bmatrix} / \$^* - [-syl]^* \begin{bmatrix} +syl \\ -stress \end{bmatrix} [-syl]^* \\ \begin{bmatrix} +syl \\ -stress \end{bmatrix})^* [-syl]^* \begin{bmatrix} +syl \\ -stress \end{bmatrix} [-syl]^* [+stress] \*$

Another example of a right to left process affecting alternate vowels is the rule postulated by Havlik (1889) to account for the development of the jers (probably to be interpreted as short high vowels) in Slavic languages. According to this rule, if successive syllables containing jers are counted from the right, then a jer in an even-numbered syllable is retained, though usually changing in quality, while a jer in an odd-numbered syllable is lost. (cf. Shevelov 1965: 452-3). In Russian the retained jers become mid vowels. For example, underlying (or historically earlier) forms *čitič* 'reader nom. sg.' and *čitica* 'reader gen. sg.' yielded *črec* and *čētica*, respectively. Havlik's rule, in its Russian version, can easily be formalized as in (53) if we are permitted left-linear application.

$$(53) \text{ (a) } L: \begin{bmatrix} +syl \\ -tns \end{bmatrix} \rightarrow \begin{bmatrix} -high \\ (\sigma) \end{bmatrix} / \$^* - [-syl]^* \begin{bmatrix} +syl \\ +hi \end{bmatrix} [-tns] \$^* \\ \text{ (b) } \begin{bmatrix} +syl \\ +hi \end{bmatrix} \rightarrow \emptyset / \$^* - \*$

An attempt to formulate (53a) in simultaneous or right-linear fashion will lead to even worse results than in the case of the Tübatulabal alternating stress rule, as can be verified by the diligent reader.

Consider too the rule of Eastern Ojibwa which turns *o* and *i* into *w* and *y*, respectively, before another vowel (Bloomfield 1956: 4-5). As Bloomfield tells us, this rule must be applied from right to left; for example, *enitioak* 'men' becomes *enitiwak*. Simultaneous or left-linear application would sometimes give the wrong result; in particular it would change *enitioak* into **enirywak*. To correct this situation we would have to complicate the rule considerably. The best we would do, it seems, is to say that *o* and *i* become *w* and *y* when occurring before an even number of *o*'s and *i*'s that are followed in turn by a vowel that is prenasal, final, or not *o* or *i*.

Certain vowel harmony rules seem also to be best formulated in linear fashion. Consider the rule of Yawelmani Yokuts which rounds a vowel that is preceded by a round vowel of the same height (Kuroda 1967: 13-15, 43-5; Kisseberth 1969). The effects of this rule are exemplified in (54)

| | | | |
|------|----------------|---------------|---------------------------------------|
| (54) | Before harmony | After harmony | Gloss |
| | hudhin | hudhun | 'recognize (aorist)' |
| | hudal | hudal | 'recognize (dubitative)' |
| | gophin | gophin | 'take care of an infant (aorist)' |
| | gopal | gopol | 'take care of an infant (dubitative)' |
| | mutmixhin | mutmuxhun | 'swear (comitative aorist)' |
| | hubuxasit | hubuxasit | 'choose (exclusive passive aorist)' |

As can be seen from the next to last example, rounding harmony is propagated as far to the right as possible. We can describe this process quite naturally by means of the right-linear rule (55)

$$(55) R: \begin{bmatrix} +syl \\ ahigh \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} +round \\ (\sigma) \end{bmatrix} / \$^* \begin{bmatrix} +syl \\ ahigh \\ +round \end{bmatrix} [-syl]^* - \*$

This formulation refers just to the nearest preceding vowel to determine whether rounding of the vowel under consideration

should take place, in the spirit of our informal description above. It should be clear that we could not adopt this approach under simultaneous or left-linear application, for each of these types of application considers only the original input form of the left-hand context. Thus we would obtain the incorrect form *mupnixhin* from *mupnixhin*. What we must do when deprived of right-linear application is allow a vowel to be rounded if it is preceded anywhere in the word by a round vowel of the same height, provided that no vowel of a different height intervenes. The necessary schema, given in (56), is much like (55) but requires an additional mention of the feature specification *zhigh*.

$$(56) \begin{bmatrix} +\text{syll} \\ \text{zhigh} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} +\text{round} \\ (\sigma) \end{bmatrix} / \$^* \begin{bmatrix} +\text{syll} \\ \text{zhigh} \\ +\text{round} \end{bmatrix} \left\{ \begin{bmatrix} -\text{syll} \\ \text{zhigh} \end{bmatrix} \right\} * - \*$

From a study of Hetzron's work (1967: 178-9, 193; 1969: 8-9, passim; and personal communication) it appears that Southern Agaw has a right-to-left rule that makes a nonlow vowel high when the next following vowel is *i*. The effect of this rule is illustrated below (tone marks omitted).

gomejanta gomejanti molegeska molegesi (input)
 multiqisi (raising rule)
 (feminine) (masculine) (plural) (singular)
 'one who is in haste' 'monk'

If my interpretation of this rule is correct, we can express it by means of the left-linear formulation (57).

$$(57) L: \begin{bmatrix} +\text{syll} \\ -\text{low} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} +\text{high} \\ (\sigma) \end{bmatrix} / \$^* - [-\text{syll}]^* \begin{bmatrix} +\text{syll} \\ +\text{high} \\ -\text{back} \end{bmatrix} \*$

A simultaneous or right-linear formulation would require us to say that a nonlow vowel is raised when followed anywhere within the word by *i*, provided that no back vowel intervened (underlying front vowels are all nonlow). We would, then need

the schema (58), which, though much like the schema in (57), mentions the feature specification —back twice instead of once.

$$(58) \begin{bmatrix} +\text{syll} \\ -\text{low} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} +\text{high} \\ (\sigma) \end{bmatrix} / \$^* - \left\{ \begin{bmatrix} -\text{syll} \\ -\text{back} \end{bmatrix} \right\}^* \begin{bmatrix} +\text{syll} \\ +\text{high} \\ -\text{back} \end{bmatrix}$$

Consonants as well as vowels may be subjected to processes of a linear character. We mention first a rather simple and commonplace example. In Russian and certain other languages an obstruent cluster becomes voiced or voiceless throughout, according as the final member of the cluster is voiced or voiceless. Thus Russian *ksgibu* 'toward the bend' is actually pronounced *gzgibu*. The left-linear formulation of the rule, given in (59) seems superior to the optimal simultaneous or right-linear formulations (60).

$$(59) L: \begin{bmatrix} -\text{son} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} \alpha\text{voice} \\ (\sigma) \end{bmatrix} / \$^* - \begin{bmatrix} -\text{son} \\ \alpha\text{voice} \end{bmatrix} \*$

$$(60) S, R: \begin{bmatrix} -\text{son} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} \alpha\text{voice} \\ (\sigma) \end{bmatrix} / \$^* - [-\text{son}]^* \begin{bmatrix} -\text{son} \\ \alpha\text{voice} \end{bmatrix} \left\{ \begin{bmatrix} +\text{son} \\ \emptyset \end{bmatrix} \right\} \*$

The schema in (60) must explicitly identify the final member of the obstruent cluster as the agent determining the voicing of the obstruent under consideration, whereas the schema in (59) needs mention only the immediately following obstruent.

Another example of a linear process affecting consonants is to be found in Tshiluba.³ Consider the verbal suffixes *-i-* (benefactive) and *-ije* (simple past). We illustrate their use first with the root *-kwar-* 'take'.

| | |
|-------------|------------------|
| kukwata | 'to take' |
| ukwačile | 'he took' |
| kukwačila | 'to take (ben.)' |
| ukwačidvile | 'he took (ben.)' |

³ Our statements about Tshiluba are based largely on work with a native informant, M. Pierre Mulumba. This work was carried out during the academic year 1969-70 at the University of California, Santa Barbara, under the auspices of the Linguistics Program. For further information about Tshiluba see Buessens (1946) and Coupez (1954).

The changes *t* to *t'* and *l* to *d'* before *i* are quite regular. Consider now the parallel paradigmatic forms of the verb -*d'ɪm*- 'cultivate':

| | |
|-------------------------|------------------------|
| kud ^ɥ ɪma | 'to cultivate' |
| ud ^ɥ imɪne | 'he cultivated' |
| kud ^ɥ imɪna | 'to cultivate (ben.)' |
| ud ^ɥ imɪɲine | 'he cultivated (ben.)' |

The rule is that *l* becomes *n* when the nearest preceding consonant is nasal, with *n* later becoming palatalized before *i*. The form *ud^ɥimɪɲine*, derived from *ud^ɥimɪɲile*, indicates that the rule proceeds from left to right. Since *n* and *l* are the only coronal sonorants in Tshiluba we can give the rule in a right-linear formulation as follows:

$$(61) \quad R: \begin{bmatrix} +\text{son} \\ +\text{cor} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} +\text{nas} \\ (\sigma) \end{bmatrix} / \$*[+\text{nas}][+\text{syll}] - \$*$$

In a simultaneous or left-linear formulation we would have to write (62).

$$(62) \quad S, L: \begin{bmatrix} +\text{son} \\ +\text{cor} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} +\text{nas} \\ (\sigma) \end{bmatrix} / \$*[+\text{nas}][(+\text{syll}) \begin{bmatrix} +\text{son} \\ +\text{cor} \end{bmatrix}]^* - \$*$$

The complication arises from the fact that we must refer to the left context in its original input form (e.g. ud^ɥimɪɲ— rather than ud^ɥimɪni—). Thus we must state explicitly that *l*'s can intervene between the *l* being considered and the preceding *n* that would trigger the nasalization. In right-linear application the intervening *l*'s will have already become *n*'s themselves, so to speak, so that the environment needn't be complicated to account for them.

The evidence considered so far favors a formalism that allows both right-linear and left-linear rules. With these types of rules available we might naturally inquire whether we need simultaneous ones as well. Apparently the answer is no, for it is difficult

to find a genuine case where simultaneous mode of application yields a better formulation than either of the linear modes. We can, however, get an idea of what such a case might be like by considering the effects of the second singular prefix in the Terena language (Bendor-Samuel 1960 and Langendoen 1968: 109-10). According to Langendoen's interpretation of Bendor-Samuel's work, this prefix has the underlying phonological shape *y*. When attached to a stem beginning with a vowel, no further change takes place. On the other hand, if the *y* is prefixed to a form beginning with a consonant, then the first non-*i* vowel undergoes a change, as follows: *u* and *e* become *i*, and *a* and *o* become *e*. Subsequently the *y* is deleted. Thus we have such forms as the following:

| | | | |
|----------|--------------|-----------------------|---------------|
| nokone | 'he needs' | nekone (<ynekone) | 'you need' |
| kurikena | 'his peanut' | kirikena (<ykurikena) | 'your peanut' |
| piho | 'he went' | pihe (<ypiho) | 'you went' |

What we have just described is the main clause of the *y*-prefix rule. There is in addition a special clause that raises an *e* to *i* in a *y*-prefixed word if all the other vowels in the word are also *e*'s. The special clause can be thought of a preceding the main clause. An example is *yxexere* 'your side', which becomes *yxiviri* by the special clause (ultimately the *y* is deleted). On the other hand, a word of the form *yBeCeDi*, where B, C, and D are consonants, would supposedly be untouched by the special clause because not all its vowels are *e*'s. The word would then pass on to the main clause, becoming *yBiCeDi*. Similarly, *yBiCeDe* would be unaffected by the special clause and would consequently become *yBiCIDE* by the main clause. It seems that only a simultaneous formulation of the special clause would be natural, since it is apparently the original input form of both the left and right contexts that determines whether an *e* will be raised. Linear application, which requires us to refer to the output form of either the left or the right context, would force us to express the special clause in a more roundabout way. Perhaps the best we could do is introduce the following two rules, each of which could be left-linear or right-linear.

- (i) Change e to E in a y -prefixed word when every vowel to the left or right is e or E ;
 (ii) Change E to i .

We assume here that E is some vowel that does not otherwise occur in Terena.

The special clause of the y -prefix rule of Terena is not in fact a clear case of an essentially simultaneous rule. There is a subtle but crucial difference between Langendoen's description of the rule and Bendor-Samuël's 1960 statement. What Bendor-Samuël actually says is that if the first one or more vowels of a word are all e 's, then these e 's are all raised to i . This implies, apparently, that we need only refer to the left context of an e to determine whether the e should be raised to i . Thus *yBeCeti* would indeed be affected, becoming *yBiCidi*. It is difficult to determine from Bendor-Samuël's data whether this is true; if it is, a left-linear formulation will be completely adequate.

Although the Terena example fails to provide us with a convincing case of simultaneity, it does present us with another type of rule, not hitherto discussed, which is interesting to take note of. Consider (63), which in its environment part is a near optimal formulation of the main clause of the Terena y -prefix rule.

$$(63) \left[\begin{array}{l} +\text{syl} \\ +\text{back} \\ \sigma \end{array} \right] \rightarrow \left[\begin{array}{l} -\text{low} \\ -\text{back} \\ (\sigma) \end{array} \right] / \left[\begin{array}{l} -\text{cons} \\ -\text{syl} \\ -\text{back} \end{array} \right] \left[\begin{array}{l} -\text{syl} \\ -\text{back} \\ (\sigma) \end{array} \right] \left\{ \begin{array}{l} -\text{syl} \\ -\text{back} \\ +\text{high} \end{array} \right\} *-\$*$$

Right-linear application cannot be associated with this schema because it would change *ykurikena* to **krikine* instead of the correct *krikena*. However, both left-linear and simultaneous application give the correct results. The reverse situation is exemplified by Dahl's law, a rule occurring in a number of Bantu languages of East Africa (Bennett 1967). The Southern Kikuyu version of this rule causes a k to become y when the next following

consonant in the word is an underlying voiceless stop. Any number of vowels may intervene between the k and the voiceless stop. As so stated, the rule can apply either simultaneously or right-linearly, producing, for example, *nyayakeroma* from *neka-kaakeroma* 'he will bite him'. (I am indebted to Leonard Talmy for this example.) A left-linear application would yield incorrectly the form **neka-yakeroma*, and consequently some other (and more complicated) means of describing the change would be necessary.

To summarize, we have found a fair number of rules that favor a right-linear formulation and a fair number that favor a left-linear formulation. In addition we have given an example of a rule that can be regarded equally well as left-linear or simultaneous, though not plausibly as right-linear, and another example that can be regarded as right-linear or simultaneous though not plausibly as left-linear. Then, of course, there are many rules (perhaps the majority) which are simply indifferent as to mode of application, working equally well in right-linear, simultaneous, and left-linear mode. The Sanskrit nasal retroflexion rule is of this kind. However, we have failed to find a convincing example of an exclusively simultaneous rule. These results suggest that we allow only right-linear and left-linear rule application in phonology.

$$(64) \quad LI: \begin{bmatrix} +sy^l \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} +stress \\ \sigma \end{bmatrix} / \$^* \begin{bmatrix} +sy^l \\ -stress \end{bmatrix} [-sy^l]^* - \$\*$

This rule will give the correct stress pattern to *natáppikkáppikkai* by virtue of the following application:

natáppikkáppikkai
 natáppikkáppikkai (by subrule a → á/nat — pikkáppikkai)
 natáppikkáppikkai (by subrule a → á/natáppikk — ppikkai)
 natáppikkáppikkai (by subrule a → á/natáppikkáppikk — i)

Notice that the elementary rule

$i \rightarrow i/natáppikkáppikkai$,

which would put a stress in the wrong place, is not a subrule of (62) even though it is subsumed under the schema in (62). The reason becomes apparent when we set

$R = natápp,$ $R' = nat$
 $P = i,$ $P' = a,$
 $S = káppikkai,$ $S' = pikkáppikkai$
 $Q = i,$ $Q' = á.$

Then the elementary rule under consideration will be $P \rightarrow Q/R - S$. However, the schema in (62) also subsumes $P' \rightarrow Q'/R' - S'$, where $RPS = R'P'S'$ and R' is shorter than R ; this situation prevents $P \rightarrow Q/R - S$ from being a subrule of (62) according to the definitions in the preceding paragraph. Note also that $á \rightarrow á/nat - pikkáppikkai$ cannot be a subrule of (62) because it is vacuous; consequently the application displayed above is complete and does not continue indefinitely.

In parallel fashion we could define a notion of right-iteration, to take the place, perhaps, of left-linear application.

Although the linear rules discussed in the preceding chapter would work as well in left- or right-iterative fashion, we have some slight reason to prefer the linear formalization. First of all, left or right iteration allows for some peculiar effects going beyond the bounds of finite-state processing, effects which we will probably

ALTERNATIVES TO LINEAR RULES

6

A right linear rule moves inexorably rightward through an input string. Having changed a segment, a right linear rule moves on, never changing that segment again nor changing anything to the left of that segment. Although this is one plausible way of formalizing left-to-right processing, it is not the only way, and it might be wrong. We might have been better advised to define a restricted iterative rule that changes only the leftmost possible segment at each step of application. Such a rule we might call left-iterative, representing it in the form $LI:X$, where X is a schema subsuming elementary rules. To achieve the effect of leftmost application formally, we might say that $P \rightarrow Q/R - S$ is a subrule of $LI:X$ if and only if X subsumes $P \rightarrow Q/R - S$ but subsumes no elementary rule $P' \rightarrow Q'/R' - S'$ where $R'P'S' = RPS$ and R' is shorter than R . We might further require that $P \neq Q$ in order to exclude vacuous subrules. We would then say that $LI:X$ maps U into V if and only if there is a sequence U_1, \dots, U_n of phonological strings such that

- (i) $U = U_1$,
- (ii) for each i , $1 \leq i \leq n - 1$, some subrule of $LI:X$ has U_i as input and U_{i+1} as output,
- (iii) U_n is not the input to any subrule of $LI:X$, and
- (iv) $U_n = V$.

The Southern Paiute stress rule, for example, could just as well have been regarded as left-iterative, having the form of (64).

want to exclude. For example, rule (65) will convert each string of the form $PCAA_1 \dots A_n B_1 \dots B_m B$, where C, B_1, \dots, B_m, B are consonants and A, A_1, \dots, A_n are vowels, into $PCAA_1 \dots A_n \text{-}m B$ or $PCAB_{n+1} \dots B_m B$ according as $n \cong m$ or $m > n$.

$$(65) \quad LI: [+syll] [-syll] \rightarrow \emptyset / \$^* [+syll] - [-syll] * [-syll]$$

The right and left iterative modes of application would have the further disadvantage of forcing a third mode of application on us. To see this, consider first any rule which switches the value of some feature, such as the first part of the English vowel shift as described by Chomsky and Halle. This rule can be given as

$$(66) \quad \begin{bmatrix} +syll \\ +tns \\ +stress \\ -low \\ zhigh \end{bmatrix} \rightarrow \begin{bmatrix} -zhigh \\ (\sigma) \end{bmatrix} / \$^* - \*$

The absence of mode designator is intended to mean that the rule can be applied in either right-linear or left-linear fashion (in fact, it can also be applied simultaneously). Either way the rule will convert $k\acute{a}w$ 'cow' into $k\acute{a}w$ (ultimately $k\acute{a}w$), by virtue of an application having the form $k \overset{f}{\delta} w$. Suppose, however, that (66) were left-iterative or right-iterative. Then it would keep on applying to $k\acute{a}w$ forever, changing it first to $k\acute{o}w$, then back to $k\acute{a}w$, then to $k\acute{o}w$ again, and so on. To make (66) work properly we would have to introduce some third type of application, the most natural choice being simultaneous.

Again, consider the rule of Chipewyan which devices a continuant consonant that immediately follows a voiceless continuant consonant (Li 1946: 400). The rule causes *tsz\acute{a}ih* 'I split' to become *tesz\acute{a}ih* (ultimately *tes\acute{a}ih*). It will also cause *n\acute{a}sz\acute{e}* 'I am hunting' to become *n\acute{a}stz\acute{e}* (ultimately *n\acute{a}sz\acute{e}*). Notice that *n\acute{a}stz\acute{e}* will not become **n\acute{a}stz\acute{e}*; it is the immediately preceding sound in the original input that determines whether devoicing of a continuant consonant will take place. The appropriate

effect can be achieved by associating simultaneous or left-linear application, though not right-linear application, with the following schema:

$$\begin{bmatrix} -syll \\ +cont \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} -voice \\ (\sigma) \end{bmatrix} / \$^* \begin{bmatrix} -voice \\ +cont \end{bmatrix} - \*$

On the other hand, both left-iterative and right-iterative application of this schema would turn *n\acute{a}sz\acute{e}* incorrectly into **n\acute{a}stz\acute{e}*. Again, some third mode of application would have to be introduced to account for this case correctly.

Some of the defects of left and right iteration could be cured as follows. We might say that a subrule $P \rightarrow R/Q - S$ changes RPS into $R < Q > S$, where P contains no angle brackets. After application of a rule is completed, angle brackets are erased. The effect, roughly, is to exclude a substring from further change if it has already been changed. By this of mechanism we could exclude at least some of the nonfinite-state effects referred to above (in particular, rule (65) would now turn $PCAA_1 \dots A_n B_1 \dots B_m B$ into $PCAA_1 \dots A_{n-1} B_2 \dots B_m B$), and we could make feature-switching rules function properly. However, the Chipewyan rule would still work improperly under both left and right iteration. To correct this situation we could attempt some further refinements of left and right iteration, but it seems pointless to do so when linear application accounts correctly for all the cases discussed.

Anderson (1968) has proposed a convention of left-to-right application that is rather different from either right-linear application or left-iteration. According to Anderson's proposal we must think not of left-to-right application of a single rule but rather of a sequence of rules. When application is to begin a position marker σ is placed immediately to the left of the first vowel in the input form. The rules, each of which contains one mention of the position marker in its environment, are then applied in sequence in the usual fashion. After the sequence is applied, the symbol σ is moved rightward until it is at the immediate left of the second vowel in the string, and the rule sequence is applied again. This

procedure is repeated as long as possible, and is therefore executed as many times as there are vowels in the input string. If we identify vowels with syllables, or somehow restrict the placing of the position marker so that it appears only to the immediate left of syllable peaks, we can refer to this mode of application as the left-to-right syllabic cycle.

Let us write a left-to-right syllabic cycle in the form $LRSy^*$: (X_1, \dots, X_n) where each X_i is a schema subsuming elementary rules. Application of the cycle might then be formalized as follows. Let P and Q be phonological strings. Then $LRSy^*$: (X_1, \dots, X_n) maps P into Q if and only if there is a sequence

$$\mathcal{S} = ((R_{1,1}, \dots, R_{1,n+1}), \dots, (R_{m,1}, \dots, R_{m,n+1}))$$

of $(n+1)$ -tuples of strings such that

- (i) P is mapped into $R_{1,1}$ by the rule
 $S: [+syl, \sigma] \rightarrow {}^o\sigma/[-syl]^* - \*
- (ii) For each $i, 1 \leq i \leq m$, and each $j, 1 \leq j \leq n$, the rule
 $S: X_j$ maps $R_{i,j}$ into $R_{i,j+1}$;
 For each $i, 1 \leq i \leq m-1$, the string $R_{i,n+1}$ is mapped into $R_{i+1,1}$ by the rule
 $S: \left. \begin{array}{l} {}^o \rightarrow \emptyset / \$^* - \$^* \\ [+syl, \sigma] \rightarrow {}^o\sigma / \$^* \text{ of } [+syl] [-syl]^* - \$^* \end{array} \right\}$
- (iv) $R_{m,n+1}$ is mapped into Q by the rule $S: {}^o \rightarrow \emptyset / \$^* - \* .

The sequence \mathcal{S} is said to be an application of $LRSy^*$: (X_1, \dots, X_n) with input P , and the $(n+1)$ -tuple $(R_{1,1}, \dots, R_{m,n+1})$ is referred to as the i th round or cycle of the application.

A very simple example of a left-to-right syllabic cycle is (67).

$$(67) \quad LRSy^*: \left(\left[\begin{array}{c} +syl \\ \sigma \end{array} \right] \rightarrow \left[\begin{array}{c} +stress \\ (\sigma) \end{array} \right] / \$^* \left[\begin{array}{c} +syl \\ -stress \end{array} \right] [-syl]^* - \$^* \right)$$

This cycle contains just one rule, whose schema is identical (except for the symbol σ) with that used in (47). In fact, (67) produces the same results as (47) and is yet another way of describing the Southern Paiute stress pattern. (67a) below is an application of

(67) to the word *tuk^wápaíyú* 'during the night'. Since (67) contains only one rule, each round of (67a) is a 2-tuple of strings.

- (67a) Round 1: (^otuk^wápaíyú, ^otuk^wápaíyú)
 Round 2: (tuk^wápaíyú, tuk^wápaíyú)
 Round 3: (tuk^wápaíyú, tuk^wápaíyú)
 Round 4: (tuk^wápaíyú, tuk^wápaíyú)
 Round 5: (tuk^wápaíyú, tuk^wápaíyú)

In this application, as in all applications of (67), the odd-numbered rounds are vacuous. Usually we will display applications of left-to-right cycles in a completely vertical manner, omitting vacuous steps. Thus (67a) would appear as follows:

^otuk^wápaíyú
 tuk^wápaíyú
 tuk^wápaíyú
 tuk^wápaíyú
 tuk^wápaíyú
 tuk^wápaíyú
 tuk^wápaíyú
 tuk^wápaíyú
 tuk^wápaíyú

Anderson does not simply propose that left-to-right syllabic cycles be made available along with other more orthodox types of application. He claims at one point that phonologies of natural languages consist entirely of simultaneous rules and left-to-right syllabic cycles. To assess the strength and validity of this claim, let us first consider left-to-right syllabic cycles from the point of view of mapping power. If we assume that a LRSy cycle can contain any rule schema that subsumes elementary rules of the form $P \rightarrow Q/R\sigma S - T$ or $P \rightarrow Q/R - S\sigma T$ and that each such schema is to be applied in simultaneous mode, then we have a rule formalism that can simulate any simultaneous rule and hence any sequence of simultaneous rules. For suppose we have a single simultaneous rule N . We can assume that N has the form $S: \{A_1 \rightarrow Q_1/X_1 - Y_1, \dots, A_n \rightarrow Q_n/X_n - Y_n\}$ where the A_i are phonological units, the Q_i are phonological strings, and the X_i and Y_i are primitive schemata subsuming phonological strings.

For each i , let X'_i and Y'_i be obtained from X_i and Y_i , respectively, by replacing each phonological unit B with the expression $(\textcircled{?}) * B(\textcircled{?})^*$. Let Z be the schema

$$\left\{ \begin{array}{l} [-\text{syll}]^* \rightarrow \$*/\$*\text{of}[\text{+syll}] [-\text{syll}]^* - \\ \$* \rightarrow \$*/\$* \left\{ \begin{array}{l} \text{of} - \\ -\$*\text{of}[\text{+syll}] \end{array} \right\} [-\text{syll}]^* \end{array} \right.$$

Then the left-to-right syllabic cycle

$$LRSY: (\{A_1 \rightarrow Q_1/X'_1 - Y'_1, \dots, A_n \rightarrow Q_n/X'_n - Y'_n\}, Z)$$

which contains just one rule, albeit a complex one, will accomplish the same task as the rule N . The bracketed expression can, of course, be eliminated in favor of a primitive schema. The idea behind this construction is that the left-to-right syllabic cycle should simply idle until it reaches the last round. On the last round, the lone rule in the cycle is executed and application terminates.

With left-to-right syllabic cycles we can also carry out certain nonfinite-state processes. Consider cycle (68) which contains just one rule.

$$(68) \quad LRSY: ([-\text{syll}] \rightarrow \emptyset / [-\text{syll}] [-\text{syll}]^* - [+\text{syll}]^* \text{of} [+\text{syll}] \$*)$$

This cycle turns a string of the form $CC_1 \dots C_n A_1 \dots A_m P$, ($n \geq 0$, $m \geq 1$), where C, C_1, \dots, C_n are consonants, A_1, \dots, A_m are vowels, and P is empty or begins with a consonant, into $CA_1 \dots A_m P$ or $CC_1 \dots C_n A_1 \dots A_m P$ according as $m \geq n$ or $m < n$. We illustrate with applications of (68) to strings *cccacaac* and *ccaac*:

| | |
|--------------------------------------|----------------------|
| cccc ^o aaac | cc ^o aaac |
| cccc ^o aaac | c ^o aaac |
| ccc ^o ca ^o aac | ca ^o aac |
| ccca ^o aac | |
| cca ^o aac | caa ^o aac |
| cca ^o aac | |

The mapping power of left-to-right syllabic cycles might be reduced if a different method of applying rule schemata were stipulated. Certain strong constraints on the form of rule schemata might also have this effect, and it is possible that we have done an injustice to Anderson's theory by not imposing such constraints. However, Anderson is silent on the matter.

Although left-to-right syllabic cycles probably have sufficient mapping capacity for phonological purposes, perhaps even too much mapping capacity, they are clearly incapable of providing linguistically satisfactory formulations in many cases. One of the difficulties arises from the fact that Anderson excludes right-to-left cycles because of his mistaken belief that right-to-left processes do not occur in the phonologies of natural languages. Thus, although the right-linear rules discussed in the preceding chapter can be routinely reformulated as left-to-right syllabic cycles (we need only place a position marker to the left of each environment dash), the left-linear rules would have to be formulated with essentially the same complex and unsatisfactory schemata that are needed in the simultaneous formulations.

Suppose we extend Anderson's formalism so as to allow right-to-left syllabic cycles as well as left-to-right ones. We would still be unable to express the right-to-left nature of the Russian rule that determines the voicing of obstruent clusters according to the final member of the cluster (rule (59) in the preceding chapter). The problem here is that the right-to-left processing takes place within each cluster, proceeding segment by segment rather than syllable by syllable. We can, however, amend Anderson's theory again so that the cycles are segmental rather than syllabic. Thus the position marker will now be placed at the beginning (or end) of the input string when application is about to begin and will be moved one segment (or, more generally, one phonological unit) to the right (or left) at the end of each round. We then will have the following two types of cycles:

- LR (left-to-right segmental)
- RL (right-to-left segmental)

All the left-to-right syllabic cycles discussed above, including the nonfinite-state cycle (68), will work properly as left-to-right segmental. The same is true of the cycles formulated by Anderson, which will be discussed below.

Anderson makes an even stronger claim than the one that we have been discussing. He believes that every phonology consists of exactly one left-to-right cycle. It should be clear by now that this claim is untenable. Note that we cannot even say that each particular language is characterized by a single direction, having either all left-to-right cycles or all right-to-left cycles. In the preceding chapter we mentioned at least two languages, Tübatulabal and Eastern Ojibwa, that had both left-to-right and right-to-left rules. Tübatulabal as a left-to-right rule of alternating vowel length and a right-to-left stress rule, while Eastern Ojibwa has a right-to-left rule making glides out prevocalic short nonlow vowels and a left-to-right stress rule.

Although most of Anderson's claims appear to be wrong, it still might be true that left-to-right (right-to-left) processing is more correctly formalized as right-to-left (left-to-right) segmental cyclic rather than left (right) linear. Let us review some Finnish evidence, presented by Anderson, which seems superficially to support this view.

Finnish has a consonant gradation rule that weakens a single or geminate stop that follows a sonorant and begins a closed syllable. The standard results of weakening are indicated below.

- p becomes v,
- t becomes d,
- k is deleted,
- a geminate stop becomes a single stop.

Thus we have forms such as the following:

¹ Among the many works about Finnish in a well-known language are Fromm and Saldemini (1956), Harms (1964), McCawley (1963, 1966), Lehtinen (1967). Our assertions about the language are based mainly on work with Miss Soitua A. Takala, who served as informant in a field methods course offered in the Linguistics Department of the University of California, Berkeley, during the academic year 1964-65.

- matto 'rug (nom.)'
- maton 'rug (gen.)'
- mato 'worm (nom.)'
- madon 'worm (gen.)'

The allative plural of *tyttö* 'girl' shows that a vowel +i cluster must be regarded as tautosyllabic for the purposes of consonant gradation:

- tyttöille (before gradation)
- tyttöille (after gradation).

There is another rule which deletes any *i* or *d* that follows a non-initial [—syll] [+syll] sequence and precedes another vowel. Some effects of this rule are seen in the partitive suffix *ta*:

- mua 'land (nom.)'
- matata 'land (part.)'
- talo 'house (nom.)'
- talota (<talota) 'house (part.)'

Suppose now we attempt to order dental-stop deletion and consonant gradation in the conventional manner. On the one hand we might decide that dental-stop deletion precedes consonant gradation, in view of the fact that this ordering yields the correct derivation for the genitive singular of *ammatti* 'occupation':

- ammatin
- ammattin (dental-stop deletion)
- ammatin (consonant gradation)

The opposite ordering would incorrectly give **ammatin*, as follows:

- ammattin
- ammatin (consonant gradation)
- *ammatin (dental-stop deletion)

On the other hand, consider the so-called collective genitive plural of *harakka* 'magpie'. If consonant gradation precedes dental-stop deletion, then the correct derivation of the form is obtained:

harakkatin
 harakkadin (consonant gradation)
 harakkain (dental-stop deletion)

The other ordering, which seems to be necessary for the genitive singular of *ammatti*, yields the wrong result for the collective genitive plural of *harakka*:

harakkatin
 harakkain (dental-stop deletion)
 *harakkain (consonant gradation)
 *harakoin (other rules).

Anderson chooses to resolve this paradox by means of a left-to-right syllabic cycle. The cycle he proposes can just as well be regarded as segmental and given roughly in the form of (69). (69a) presupposes a rule that stresses the first and only the first vowel of a word. Just how the portion of (69b) to the left of the slash is to be adequately expressed is not clear; this matter will be touched on later.

(69) LR:

(a) Dental-stop deletion.

$$\left[\begin{array}{c} -\text{cont} \\ +\text{cor} \end{array} \right] \rightarrow \emptyset / \$* [-\text{syll}] \left[\begin{array}{c} +\text{syll} \\ -\text{stress} \end{array} \right] -[\sigma + \text{syll}] \$*$$

(b) Consonant gradation.

$$\left\{ \begin{array}{l} t \rightarrow t \\ t \rightarrow d \\ \cdot \\ \cdot \end{array} \right\} / \$* [\sigma + \text{son}] -[\sigma + \text{syll}] \left\{ \begin{array}{l} i \\ \emptyset \end{array} \right\} [-\text{syll}] \left\{ \begin{array}{l} [-\text{syll}] \$* \\ \emptyset \end{array} \right\}$$

The correct form of the words considered above will now result, as can be seen from the following partial derivations:

(i) ammatt^oin (69b)
 ammat^oin
 ammatⁱn

(ii) harakk^oatin
 harakk^otin
 harakkat^oin
 harakka^on (69a)
 harakkaⁱn

Though Anderson's artifice is clever, the need for it is highly questionable. Notice, first of all, that conventional rule ordering can be imposed if we split the gradation rule into at least two stages. The first stage would be much like the gradation rule in (69), except that it would merely weaken one half of a geminate stop rather than eliminating that half altogether. In other words, the gradation rule proper would change *ti* into *ti*, say, where *t* is some weakened variety of *t*, and a later rule following dental-stop deletion, would reduce *ti* to *t*. We would then have derivations such as the following:

ammatin harakkatin
 ammatin harakkadin (consonant gradation)
 ammatin harakkain (dental-stop deletion)
 ammatin harakkain (weak geminate reduction)

Despite Anderson's assertion to the contrary, there is good independent evidence that geminates pass through an intermediate stage when they are subject to gradation. To see that this is so we will have to consider a number of rules which appear at first glance to be irrelevant to the present problem.

First, although the usual results of weakening *p, t, k* in the gradation environment are *v, d,* and \emptyset , respectively, special circumstances will call forth a different output. For example, if the sonorant preceding the stop is a homorganic nonsyllabic, the stop assimilates completely to the sonorant; for example, underlying *lukenton* (genitive singular of *lukento* (> *luento*) 'lecture') becomes *luennon*. Another fact, of particular interest here, is that *k* becomes *j* in the gradation environment when the preceding sonorant is *r, l,* or *h* and the following vowel is *e*. Thus form the verb stem *särke-* 'break' we have first person singular present *säryjen* (cf.

third person plural present *särkevät*). It should be borne in mind that it is only the vowel *e* which triggers the change *k* to *j*; even the closely related vowel *i* will not do it. Thus when the first singular suffix *n* is appended to the verb stem *pyrki-* 'strive for' the phonetic form *pyrin* results according to the general rule, not *pyrin*.

Consider next the rule that deletes a short unrounded vowel when a suffixal *i* follows. The effects of this rule can be illustrated by the verb stems *pyrki-*, cited above, *hake-* 'look for', and *ottaa* 'take'. If the third person suffix *vät* alone is added to these stems, the resulting forms are simply *pyrkivät* 'they strive for' (with vowel harmony in the suffix), and *hakevat* 'they look for', *ottavat* 'they take'. If the past tense suffix *i* is inserted too, giving the underlying strings *pyrkivät*, *hakeivat*, and *ottivat*, the phonetic results will be *pyrkivät*, *hakeivat*, and *ottivat*. If we add the past and first singular suffixes to the stems in question, obtaining the underlying forms *pyrkiiin*, and *hakiin*, and *otitiin*, both gradation and unround-vowel-deletion will apply, yielding phonetic *pyrin*, *hain*, and *otin*.

A verb stem ending in *rke*, *lke*, or *hke*, will of course also be subject to unround-vowel-deletion. The crucial question is what happens when both unround-vowel-deletion and gradation apply. Applied in the order just stated, the *k* would be deleted; thus the first singular past of *särke-* would be obtained by the following derivation:

särkein
särkin (unround-vowel-deletion)
särin (gradation)

The output here, however, is incorrect; the form actually used is *särjin*. Consequently, consonant gradation must precede unround-vowel-deletion, and we must have the derivation

särkein
särjein (gradation)
särjin (unround-vowel-deletion)

Another rule we must consider is the one that changes noninitial *t* to *s* before *t*. This "dental-stop sibilant" rule, together with a

rule that raises final *e* to *i*, accounts for nominative singular *käsi* 'hand' in face of the essive singular *käsenä*, genitive *käden* (with gradation), and so on. It also accounts for the past tense forms of the many verb stems that end in *t*. An example is *haluta-* 'want' with third singular imperative *halukoon*, infinitive *haluta* (from *haluttah* by consonant gradation and final *h* deletion), first singular past *halusin* (< *halutin* by sibilant). The last form indicates that dental-stop deletion must follow sibilant; otherwise we would obtain **halutin*. The stem *kipet-* 'climb' behaves in a similar manner; we have *kiivetkään* (< *kipetkoon*), *kiivetä* (< *kipettah*), *kipesin* (< *kipetin*).

Consider now a verb stem like *tunte-*, 'feel' infinitive *tuntea* from underlying *tuntetah*. In the present we have first singular *tunen*, from *tunteti* by consonant gradation, and third plural *tuntevat*. In the past we have first person singular *tunsin* and third plural *tunsivat*, from underlying *tuntein* and *tunteivat*. Clearly, unround-vowel-deletion must precede sibilant of *t* if these forms are to be derived correctly.

The data we have considered so far imposes the following order on the rules under discussion:

consonant gradation
 unround-vowel-deletion
 sibilant of dental stops
 dental-stop deletion

If we accept this ordering, which seems to be unavoidable even in a left-to-right cycle, we must perforce accept an intermediate stage in the gradation of geminate stops. For if the gradation rule immediately reduced *tt* to *t* then the rules given above would turn *ammattiin* into **ammasin* rather than the correct form *ammatin*. Gradation must, then, weaken *tt* to some intermediate form which cannot be reduced to *t* until after both dental-stop sibilant and dental-stop deletion have applied. The precise nature of this intermediate stage is not altogether clear, to be sure; its historical analogue is denoted *ti* in the handbooks (cf. Fromm and Sadeniemi 1956: 35-6) in accordance with a hypothesis that the first part

of the cluster was weakened (laxed?). Nevertheless, the basic point remains. If geminate stops must pass through an intermediate stage anyway when subject to gradation, a major motivation for the left-to-right cycle vanishes.

Note that geminate stops are not the only clusters that must pass through an intermediate stage under the rule ordering that seems to be necessary. Consider again the first singular present and past of *tunte*-. We saw above that these forms had the underlying representations *tuntēn* and *tuntein*. The phonetic realization of *tuntēn* is *tunnen*, which arises straightforwardly by gradation.

The phonetic realization of *tuntein*, however, is *tunsiin*, which must have undergone dental-stop sibilation. But if gradation changes *n* immediately to *m*, our rules give the following incorrect derivation:

tuntein
 tunnein (gradation)
 *tunnin (unround-vowel deletion)
 (no further rules applicable)

Apparently, then, we must assume that a single stop in gradation environment is first weakened by some simple operation such as laxing or voicing, regardless of what the preceding sonorant is. This weakened stop is still subject to sibilation if it is dental; if it escapes sibilation it is later assimilated to a preceding homorganic consonant, if any. If a weakened stop escapes all these rules it becomes *v*, *d*, or \emptyset depending on whether it is labial, dental, or velar. Thus if we denote the weakened stops by the noncommittal symbols β , *t*, *k*, we have derivations something like the following:

tunten tuntein
 tunten tunnein (gradation)
 tunfin tunnin (unround-vowel deletion)
 tunsin (sibilation)
 tunnen (assimilation)

Another problem that has concerned Anderson is the formation of the illative case of Finnish nouns and adjectives. In order to

understand what this is about we must consider some further rules.

Stems ending in a consonant add a linking vowel *-e-* before most suffixes that begin with a consonant. Compare the following forms

talo 'house (nom.)' sisar 'sister (nom.)'
 talon 'house (gen.)' sisaren 'sister (gen.)'

In both cases we assume that the genitive is formed by adding *n* to the noun stem; *sisar*+*n* then becomes *sisaren* by the *e*-insertion rule.

E-insertion does not take place before certain suffixes, however; among these exceptional suffixes is the partitive *ta*. Thus we have *talota* 'house (part.)', from *talota* by dental-stop deletion, and *sisarta* 'sister (part.)'. Whether the difference between the two types of suffixes can or ought to be described in phonological terms or in terms of diacritic features on morphemes is a question that will not concern us here. We will assume that the linking *e* appears in its proper place prior to the application of any of the rules to be discussed (in fact, prior to any of the rules discussed above).

Consider now some noun and adjective stems that end in *s*. We have *mies* 'man', *kirves* 'axe', *vieras* 'foreign', and *uros* 'male', with genitives and partitives in the singular as follows:

genitive: miehen kirveen vieraan uroon
 partitive: miestä kirvestä vierasta urosta

Anderson posits three rules that will account for the above genitives. These rules, which are motivated by other evidence as well and which we see no strong reason to question, are essentially the following:

- (i) intervocalic *s* after a nonfirst vowel becomes *h*
 ("s-weakening") (This rule has numerous exceptions,
 e.g. *kisä* 'pyrites'),
 (ii) *A*/*h*e, where *A* is any vowel, becomes *AhA* ("vowel-
 copying"),

- (iii) intervocalic *h* after a noninitial [—syl][+syl] sequence is deleted.

The derivations of the genitives are, then,

| | | | | |
|------------------------|--------------------------|------------------------|--------------------------|------------------------|
| <i>mies</i> + <i>n</i> | <i>kirves</i> + <i>n</i> | <i>uros</i> + <i>n</i> | <i>vieras</i> + <i>n</i> | (<i>e</i> -insertion) |
| <i>miesen</i> | <i>kirvesen</i> | <i>urosen</i> | <i>vierasen</i> | |
| <i>miehen</i> | <i>kirvehen</i> | <i>urohen</i> | <i>vierahen</i> | (<i>s</i> -weakening) |
| | | <i>urohon</i> | <i>vierahan</i> | (vowel-copying) |
| | <i>kirveen</i> | <i>uroon</i> | <i>vieraan</i> | (<i>h</i> -deletion) |

In *miehen* the *h*, though intervocalic, is not deleted because it is preceded by a diphthong. (A more refined analysis would reveal that the underlying form of *mies* is *mies_h*, so that it would in fact be a long vowel that inhibits *h*-deletion.)

Let us turn now to the illative. We consider the nouns *maa* 'land', *puu* 'tree', *talo* 'house', *koira* 'dog', and *katu* 'street', in the illative singular and plural.

| | | | | | |
|-----------|---------------|---------------|-----------------|----------------|-----------------|
| Singular: | <i>maahan</i> | <i>puuhun</i> | <i>taloon</i> | <i>koiraan</i> | <i>katuun</i> |
| Plural: | <i>maihin</i> | <i>puihin</i> | <i>taloihin</i> | <i>koiriin</i> | <i>katuihin</i> |

We account for all these forms with the rules we have plus one additional rule, provided we assume with Anderson that the underlying form of the illative suffix is *sen*. The additional rule, given in a somewhat different version by Anderson, deletes the middle vowel of a triphthong ("triphthong reduction"). Four derivations should make the matter clear

| | | | | |
|---------------|---------------|-----------------|--------------------------|------------------------|
| <i>maasen</i> | <i>maisen</i> | <i>koirasen</i> | <i>koiraisen</i> | |
| | | <i>koirisen</i> | (unround-vowel deletion) | |
| <i>maahen</i> | <i>maahen</i> | <i>koirahen</i> | <i>koirihen</i> | (<i>s</i> -weakening) |
| <i>maahan</i> | <i>maahin</i> | <i>koirahan</i> | <i>koirihin</i> | (vowel-copying) |
| | | <i>koiraan</i> | <i>koiriin</i> | (<i>h</i> -deletion) |
| | <i>maihin</i> | | (triphthong red.) | |

The rules just given predict the illative form correctly when the underlying stem consists of one syllable or ends in a short vowel or a nonspirantal consonant. We must assume that illative *sen* is one of those suffixes that requires *e*-insertion; in this way we

obtain the correct illative singulars *miehen* (< *miesesen*) and *sisareen* (< *sisaresen*).

Let's turn now to the disyllabic stems ending in *s* which we mentioned above. The illative singulars *kirves*, *uros*, and *vieras* are

| | | |
|------------------|-----------------|------------------|
| <i>kirveseen</i> | <i>urooseen</i> | <i>vieraseen</i> |
|------------------|-----------------|------------------|

There are two alternatives in the plural:

| | | |
|-----------------------|-----------------|-------------------|
| (a) <i>kirveisiin</i> | <i>uroisiin</i> | <i>vieraisiin</i> |
| (b) <i>kirvehin</i> | <i>uroihin</i> | <i>vieraihin</i> |

There are some other disyllabic stems which in the nominative singular end in a short vowel but behave in other respects like stems ending in *s*. Anderson assumes they end in a final *h*, which is deleted when no suffix follows as well as in the environment of the *h*-deletion rule introduced previously. *h* followed by a consonant results in gemination of that consonant. Thus from the stem *kiiruh* 'hurry' we have the following case forms:

| | |
|-----------|------------------------------|
| nom. sg. | <i>kiiru</i> |
| part. sg. | <i>kiirutta</i> |
| gen. sg. | <i>kiiraun</i> |
| ill. sg. | <i>kiiruuseen</i> |
| ill. pl. | <i>kiiruisiin, kiiruihin</i> |

Polysyllabic stems ending apparently in a long vowel behave in many respects like these ending in a spirant. From the stem *tieno* 'neighborhood' we have

| | |
|-----------|------------------------------|
| nom. sg. | <i>tienoo</i> |
| part. sg. | <i>tienootta</i> |
| gen. sg. | <i>tienoon</i> |
| ill. sg. | <i>tienooseen</i> |
| ill. pl. | <i>tienoisiin, tienoihin</i> |

The nonillative case forms of polysyllabic stems ending in a spirant or long vowel can be accounted for by the rules already introduced, if appropriately ordered in the conventional manner. The illatives,

with their peculiar vowel-lengthening and retention of *s*, offer some difficulty. Anderson accounts for them with a left-to-right syllabic cycle, which we give as a roughly formulated segmental cycle in (70) below.

(70) LR: (a) Triphthong reduction.

$$[+syll] \rightarrow \emptyset / \$* [+syll] - [+syll] [-syll] [e [+syll]] \$*$$

(b) Illative lengthening.

$$\begin{bmatrix} +ill \\ \sigma \end{bmatrix} \rightarrow \sigma \sigma / \$* \begin{bmatrix} +syll \\ -stress \end{bmatrix} \begin{matrix} \tau \\ \tau \end{matrix} \left\{ \begin{matrix} i \\ \emptyset \end{matrix} \right\} [-syll] o - \$*$$

(c) S-weakening.

$$s \rightarrow h / \$* \begin{bmatrix} +syll \\ -stress \end{bmatrix} [-e [+syll] [-syll]] \$*$$

(d) Vowel-copying.

$$e \rightarrow \sigma / \$* \begin{bmatrix} +syll \\ -stress \end{bmatrix} h^o - \$*$$

(e) H-vocalization.

$$h \rightarrow \sigma / \$* [-syll] \begin{bmatrix} +syll \\ -stress \end{bmatrix} [-e [+syll]] \$*$$

Of the rules in this cycle, triphthong reduction, *s*-weakening, and vowel-copying have been already introduced above, and *h*-vocalization takes the place of *h*-deletion. Illative lengthening is a new rule applying to one idiosyncratic case morpheme. Instead making use of the diacritic feature [illative], presumably introduced by a general convention that distributes syntactic class information to the individual segments of a morpheme, we could have regarded illative lengthening as a 'minor rule' from which all morphemes except the illative are exempt.

The illative singular and plural of *uros* can now be derived in part as follows:

| | | |
|------------------------|------------------------|-------|
| uros ^o esen | uros ^o isin | |
| uroh ^o esen | uroh ^o isin | (70c) |

| | |
|------------------------|-------|
| uroh ^o osen | (70d) |
| uroo ^o osen | (70e) |
| uroo ^o isen | |
| urooos ^o en | |
| uroos ^o en | (70a) |
| uroos ^o in | (70b) |
| uroos ^o een | (70c) |

The alternative illative plural arises from a different ordering of (70a) and (70b); for example:

| | |
|-------------------------|-------|
| uros ^o isin | (70c) |
| uroh ^o isin | (70e) |
| uroo ^o isin | |
| urooi ^o isin | |
| uroois ^o in | (70b) |
| uroois ^o in | (70a) |

Notice that in these derivations the triphthong reduction rule comes into play only when the position marker has moved to the immediate left of the first vowel that follows the triphthong. This feature of Anderson's solution is crucial in the derivation of the illative plural form *uroisin*, for it allows illative lengthening, formulated with *o* before the environment dash, to take place prior to triphthong reduction despite the fact that the illative vowel is to the right of the triphthong. Thus, when illative lengthening is ordered before triphthong reduction, *uroois^oin* will first become *uroois^oin* and then *uroois^oin*. Then *s*-weakening, formulated so as to take place only before a short vowel, is blocked; thus the derivation of the nonexistent form **uroihin* is prevented and the desired form *uroisin* is generated.

Unfortunately, there is a serious difficulty with Anderson's formulation of the triphthong reduction rule. Consider the verb stem *saz-* 'receive'. The infinitive and first singular present are *sazda* (from *satah* by consonant gradation and final *h* deletion) and *saan*. Now the past, formed by adding *i* plus the personal

suffixes, is always subject to triphthong reduction; thus we have first singular past *sain* from underlying *sain*. Similarly, the third singular past, where the personal suffix is \emptyset , has the form *sai*. Since no vowel occurs after the triphthong *ai* in the underlying forms of these words, Anderson's rule could not apply and *ai* would remain unreduced. Similar remarks apply to verb stems like *jo-* 'drink' and *sö-* 'eat', which have first singular presents *juon* and *syön* (from *joon* and *söön* by a vowel breaking rule) and first singular pasts *join* and *söin*, derived from *jooin* and *söiin* by triphthong reduction. A more correct and more natural formulation of triphthong reduction would seem to be

$$[+syl] \rightarrow \emptyset / \$* [+syl] - [+syl] \$*$$

With this formulation, of course, it is impossible to use the position marker exactly as Anderson does, although it might be possible to restore the essence of Anderson's solution by rejinggling the position marker in this and other rules. Instead of showing how this could be done, however, we will proceed to some further difficulties not as easily resolved.

The rule of *s*-weakening was originally formulated to take place regardless of whether the following vowel is long or short, but was restricted to position before short vowel in order to account for illatives. Thus *uroosseen* (or *urooseen*) and *urootsiin* are correctly prevented from undergoing *s*-weakening, as explained above. However, there seems to be no other motivation for the short vowel restriction, although indeed there seems to be no counter-evidence to it either. On the other hand, there are many apparent exceptions to *s*-weakening which are not explained by the restriction. Thus we have *kisii* 'pyrites', *nielaisen* 'I swallow' (cf. *nielaisia* 'to swallow'). Whatever the explanation of these intervocalic *s*'s may be, it has nothing to do with the length of the following vowel. This fact raises the suspicion that the simpler, less restricted *s*-weakening rule might just as well be used and that the *s* illatives are to be explained in a different way.

The third and perhaps most serious difficulty with Anderson's solution is this. Notice that Anderson's derivations of the illative

plural forms of *uros* begin not with the presumed underlying form *urosien* but with *urosisin*. Anderson does not explain the reason for this, nor does he give any rule that would change *urosien* into *urosisin*. The change looks suspiciously like a case of vowel-copying, and can indeed be so regarded if we assume that *urosien* first becomes *uroithen* and that the *r* is re-introduced by a revised version of the illative lengthening rule. This approach (similar to that followed by McCawley 1966) also allows us to use the simpler versions of both the triphthong reduction rule and the rule of *s*-weakening, discussed above. Furthermore, it makes it possible to apply the rules in conventional sequence. Our reformulation, still somewhat informal because not stated entirely in feature terms, is given in (71).

(71) (a) *S*-weakening.

$$s \rightarrow h / \$* \left[\begin{array}{l} +syl \\ -stress \end{array} \right] - [+syl] \$*$$

(b) Vowel-copying

$$L: e \rightarrow \sigma / \$* \left[\begin{array}{l} +syl \\ -stress \end{array} \right] h - \$*$$

(c) H-vocalization.

$$R: h \rightarrow \sigma / \$* \left[\begin{array}{l} +syl \\ -stress \end{array} \right] - [+syl] \$*$$

(d) Triphthong reduction.

$$[+syl] \rightarrow \emptyset / \$* [+syl] - [+syl] \$*$$

(e) Illative adjustment.

$$[+ill] \rightarrow \sigma / \$* \left[\begin{array}{l} +syl \\ -stress \end{array} \right] \tau \left\{ \begin{array}{l} i \\ \emptyset \end{array} \right\} - \sigma \$$$

These rules apply in the order given, except that (71d) and (71e) can also be applied in the reverse order.

The derivations of the illative forms of *uros* now proceed as follows:

| | | |
|-----------|-----------|-------|
| urosescen | urosisen | |
| urohethen | urohithen | (71a) |
| urohohen | urohithin | (71b) |
| urooohen | urooithin | (71c) |
| urooohen | urooithin | (71d) |
| uroosecen | urooithin | (71e) |

As they stand, our rules account correctly for the illatives of stems like *uros*, which end in a spirant, but not for the illatives of stems that end in a long vowel. For example, although our rules correctly predict the illative plural forms of *tieno*, they give the wrong illative singular. By our rules, the underlying form *tienoosen* first becomes *tienoohen* by s-weakening and then *tienoohon* by vowel-copying. The illative adjustment rule then gives the incorrect result **tienoosoon*. However, we can avoid this consequence by giving *tieno* the underlying representation *tienose* or *tienohe*. The illative singular can be then derived from underlying *tienosescen* in the same way that the illative singular of *uros* is derived from *urosescen*. If we assume that unround-vowel deletion precedes s-weakening we can derive the alternative illative plurals of *tieno* from the underlying representation *tienoeseisen* as follows. First, *tienoeseisen* becomes *tienoseisen*, and then *tienoeseisen* follows the same path as *urosescen*. The nominative singular is derived straightforwardly as follows:

| | |
|---------|-------|
| tienose | |
| tienohe | (71a) |
| tienohe | (71b) |
| tienooh | (71c) |
| tienooh | (71d) |

Words like *leikkau* 'harvest' and *vapaa* 'free', which behave just like *tieno*, can be treated similarly. They would have underlying forms like *leikkase* or *leikkaihe* and *vapase* or *vapaihe*.

An objection that might be raised against our treatment of words like *tieno* is that a perfectly easy way of accounting for the partitive singular *tienoita* must be abandoned. If we assume with

Anderson that underlying form of *tieno* is *tieno*, then the retention of *t* in the partitive suffix is explained by the fact that a long vowel precedes the *t*. (cf. the environment of the cyclic dental-stop deletion rule (69a)). If we assume that the underlying representation of *tieno* is *tienose*, then when we add the partitive suffix we have the underlying string *tienoasetä*. Here the *t* is preceded by a cluster subsumed under [-syll] [+syll, -stress], which ought to trigger deletion of the *t*. The solution that immediately suggests itself is that we order dental-stop deletion after *h*-vocalization. Then we would have the derivation

| | |
|-----------|-------------------------------|
| tienoseta | |
| tienoheta | (s-weakening (71a)) |
| tienohota | (vowel-copying (71b)) |
| tienoota | (h-vocalization (71c)) |
| tienoota | (triphthong reduction (71d)) |
| | (no further rules applicable) |

The ordering of *h*-vocalization prior to dental-stop deletion seems to be supported independently by the illative plurals of noun and adjective stems ending in a spirant. For example, *uroita*, the illative plural of *uros*, can be derived properly from the underlying string *uroisita* only if *h*-vocalization precedes dental-stop deletion.

Consider, however, the noun *airut* 'herald'. The illative plural will be underlying *airutisen*. Our rules as they now stand will convert this representation into *airuithin*. However, the correct form is *airuithin*. Observe that *airuithin* would indeed be generated if, contrary to our conclusions above, we ordered dental-stop deletion before *h*-vocalization. Then we would have the derivation

| | |
|-----------|------------------------|
| airutisen | |
| airuithen | (s-weakening (71a)) |
| airuithin | (vowel-copying (71b)) |
| airuithin | (dental-stop deletion) |

h-vocalization would not apply to the last line above because the *h* is preceded by a diphthong.

The apparent conflict between forms like *urolia* and forms like *airiuhin* can easily be resolved if *h*-vocalization and dental-stop deletion are regarded as subcases of a single right-linear rule. We might have thought to combine the two processes into a single rule anyway, since their environments look very similar the way we have formulated them. The similarity is not as close as it appears, however, for we have somewhat oversimplified the environment of *h*-vocalization. Apparently, *h*-vocalization takes place after any unstressed vowel sequence that does not end in a long vowel or in vowel +*i*, whereas dental-stop deletion is restricted to position after consonant + short unstressed vowel. Some effects of this difference are illustrated below.

| | | |
|-------------|------------|----------|
| nominative: | autio | hertua |
| illative: | autioon | hertuaan |
| partitive: | autiota | hertuata |
| | 'desolate' | 'duke' |

Compare also ill. sg. *airveen* (< *airutesen*) with ill. pl. *airiuhin*. The question now is how to express gracefully the environment of *h*-vocalization. I don't know a definitive answer, but perhaps we can achieve the right effect if we take the liberty of using the following somewhat controversial devices:

- (i) brackets meaning set intersection, as described in Chapter 3;
- (ii) hyphen meaning negation (that is, -(X) subsumes everything that X doesn't).

Then the rule of *h*-vocalization and the rule of dental-stop deletion can be combined into the single right-linear rule (72).

$$(72) \quad R: \left[\begin{array}{l} [-\text{cons}] \\ [-\text{voice}] \end{array} \rightarrow \sigma / \left[\begin{array}{l} +\text{syll} \\ -\text{stress} \end{array} \right], -(\$* \left[\begin{array}{l} +\text{syll} \\ \tau \end{array} \right] \left[\begin{array}{l} i \\ \tau \end{array} \right] \right) \left. \vphantom{\begin{array}{l} [-\text{cons}] \\ [-\text{voice}] \end{array}} \right\} -[+\text{syll}] \$* \\ \left[\begin{array}{l} [-\text{cont}] \\ [+ \text{cor}] \end{array} \right] \rightarrow \emptyset / \$* \left[\begin{array}{l} +\text{syll} \\ -\text{stress} \end{array} \right]$$

It is possible that this rule can and should be further reduced by appropriate use of angle brackets or indexed braces in order to express what similarity there does exist between the left environments of the subcases.

We emphasize that our solutions to the problems under discussion are far from definitive. Numerous unresolved questions remain, but these can be answered only by a far deeper study of Finnish phonology than would be appropriate here. We could claim only that the solutions presented here are at least as plausible as Anderson's and in some respects superior, and that there is little reason to believe, therefore, that a left-to-right cycle is necessary. It is true that at least one rule, (72), must be applied in a left-to-right manner, but a right-linear formulation appears to be adequate.

7

FEATURES WITH INTEGRAL COEFFICIENTS

It appears from Cole's description (1955) that Tswana has the following system of underlying vowels:

| | | |
|---|---|------------|
| i | u | High |
| e | o | Higher mid |
| ɛ | ɔ | Lower mid |
| a | | Low |

Using the vowel features proposed in Wang (1968) we would presumably attribute the following analysis to these vowels.

| | | | | | | | |
|---------|---|---|---|---|---|---|---|
| | i | u | e | o | ɛ | ɔ | a |
| High | + | + | + | + | - | - | - |
| Mid | - | - | + | + | + | + | - |
| Palatal | + | - | + | - | + | - | - |

As contrasted with the vowel features of Chomsky and Halle (1968), hereafter referred to as SPE, the above analysis has two distinct advantages in the present example: (1) it allows for the representation of four vowel heights; and (2) the vowels specified [+mid] constitute a natural class that is opposed to the remaining three vowels, as will be seen below.

In certain environments the mid vowels are raised slightly, but not enough to reach the height of the next higher vowel; in other words *e* is raised but not to the extent of becoming *ɛ*, and *ɛ* is raised but not to the extent of becoming *i*. Phonetically, then, the following vowels are distinguished:

| | | |
|----|----|-------------------|
| i | u | High |
| e~ | o~ | Raised higher mid |
| e | o | Higher mid |
| ɛ~ | ɔ~ | Raised lower mid |
| ɛ | ɔ | Lower mid |
| a | | Low |

One of the raising environments is before the velar nasal ŋ. Thus we have *o~ŋ*, *e~ŋ*, *ɔ~ŋ*, and *ɛ~ŋ* from underlying *oŋ*, *eŋ*, *ɔŋ*, and *ɛŋ*, respectively. This rule is followed by another rule which, working from right to left, raises a mid vowel by one degree if the next following vowel is higher. We illustrate as follows:

| | | | |
|--------------|------------------------------|-----------------------------|-------------------------------|
| moxakolodi | moselesela | moseleselen | (raising before ŋ) |
| | | moselesela~ŋ | |
| moxako~lo~di | | moser~le~se~le~ŋ | (raising before higher vowel) |
| 'advisor' | 'Dichrostachys glomerata' | (locative of moselesela) | |

The question now arises as to how we are to represent the six phonetic vowel heights in feature terms and how we are to formulate the raising rules. Clearly, the four-way height distinction that is adequate for the underlying representations is insufficiently refined for the phonetic realizations. We can, however, easily account for this situation if we allow that feature coefficients on the phonetic level be given as integers rather than as pluses and minuses. When provided with such coefficients a feature is thought of as indicating the degree to which a certain phonetic quality is possessed. Thus phonetic vowel height might be thought of as given by one of the first *n* positive integers, 1 indicating the least tongue height, and *n* indicating the greatest tongue height. The scale of height given by the integers 1 through *n* should presumably be universal, defined once and for all in phonological theory. In

the absence of a generally accepted universal scale, we will assume ad hoc that $n = 6$, since this is adequate for Tswana. The rules of Tswana must at some point convert the plus and minus specifications of the high and mid features into integer specifications of the phonetic feature of height, and the raising rules will be formulated in terms of those integer specifications.

In order to handle integer specifications we will introduce a few new devices. We regard a positive integer n as a sequence of n 1's. Thus $1 = 1$, $2 = 11$, $3 = 111$, etc. The Greek letters ι , κ , λ , ... will be regarded as variables ranging over strings of 1's. We will furthermore allow expressions of the form $+IF$, where I is a string of 1's. Such an expression subsumes any phonological unit which contains the specification $+IF$ or $1IF$ where J is a string of 1's. Thus in particular, $+F$ subsumes units containing the specifications $+F$, F , $1F$, $11F$, $111F$, etc.

Suppose now that there is a rule in Tswana that converts binary height specifications into integer specifications as follows:

| | | |
|--------------|---------|----------------|
| i , | u | become 6 high |
| e , | o | become 4 high |
| ϵ , | \circ | become 2 high |
| a | | becomes 1 high |

Then the rule that raises a mid vowel before a vowel of greater height can be given as in (75):

$$(75) \quad L: \begin{bmatrix} +\text{sy}l \\ \iota 11 \text{ high} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} \iota 111 \text{ high} \\ (\sigma) \end{bmatrix} / \$* - [-\text{sy}l] * \begin{bmatrix} +\iota 111 \text{ high} \\ +\text{sy}l \end{bmatrix} \$*$$

The vowel height feature, as we have viewed it, has finitely many possibly integer coefficients. As long as this is true of all features the alphabet of phonological units remains finite, and the extended formalism allowing integers and integer variables does not increase the mapping power of phonological rules. It is otherwise with the prosodic features now to be discussed.

According to Schachter and Fromkin (1968: 106-9) the Akan

languages have an underlying distinction between high-level and low-level tone. These tones are manifested phonetically according to the following downdrift rule (after certain other rules have applied). A high tone is always on a higher pitch than an immediately neighboring low tone; however, the pitch interval between a low tone and an immediately following high tone is less than the pitch interval between a high tone and an immediately following low tone. To quantify these relations Schachter and Fromkin assume that there is a pitch interval of two degrees between a low tone and a following high tone and that there is a pitch interval of three degrees between a high tone and a following low tone. We can express this analysis in feature terms if we suppose there is an integrally valued feature of pitch, 1 pitch being the highest pitch, 2 pitch the next highest, and so on. Then the effect of the rule can be illustrated by the phrase

òbék's kùmbéé ándópá yí
s 1 1 4 2 2 2 5 3 3

Here high tone is indicated by an acute accent, low tone by a grave accent. High tone can be thought of as represented by the specification $+tone$, low tone by the specification $-tone$. The downdrift rule can be given by a series of three rules, of which the last is right-linear.

$$(76) \quad \begin{array}{l} (a) \quad \begin{bmatrix} +\text{tone} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} 1 \text{ pitch} \\ (\sigma) \end{bmatrix} / \$* - \$* \\ (b) \quad \begin{bmatrix} +\text{sy}l \\ +\text{tone} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} \iota 111 \text{ pitch} \\ (\sigma) \end{bmatrix} / \$* - \$* \\ (c) \quad R: \left\{ \begin{array}{l} \begin{bmatrix} +\text{tone} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} \iota \text{pitch} \\ (\sigma) \end{bmatrix} / \$* \begin{bmatrix} +\text{sy}l \\ -\text{tone} \\ \iota 111 \text{ pitch} \end{bmatrix} \left[\begin{array}{l} -\text{sy}l \\ +\text{tone} \end{array} \right] * - \$* \\ \begin{bmatrix} +\text{sy}l \\ -\text{tone} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} \iota 111 \text{ pitch} \\ (\sigma) \end{bmatrix} / \$* \begin{bmatrix} +\text{tone} \\ \iota \text{pitch} \end{bmatrix} [-\text{tone}] * - \$* \end{array} \right. \end{array}$$

We assume here that only syllabics can have the specification +tone, nonsyllabics being all -tone.

There are now infinitely many phonological units, since in principle [*n* pitch] segments exist for each *n*. In order to continue thinking in terms of a finite alphabet of primitive symbols we can, of course, think of a phonological string as composed of brackets, commas, pluses, minuses, 1's, and feature names, these being the formal building blocks of phonological units. We will then be in a position to consider whether the extended formalism allowing integer coefficients and variables is more powerful than the formalism of Chapters 4-5. In fact, it is more powerful, for many rules can be formulated that exceed the power of finite transducers. The Akan downdrift rule (76c), for example, is not a finite state device. To see this we can reason as follows. Let us say that *Q* is in the output set of a mapping device if and only if that device maps at least one string into *Q*. It is well known that the output set of a finite transducer is regular. Therefore, if *M* is some finite transducer that purportedly effects the same mapping as rule (76c), then the output set of both *M* and rule (76c) must be regular. However, the output set of (76c) is not regular because of the dependencies that hold between the coefficients of the pitch features within a phrase. If we regard these coefficients as strings of 1's we are obliged to say that in any output string of the form [L +tone, I pitch K -tone, J pitch H], where *K* contains no occurrence of the expression +tone and *I* and *J* are strings of 1's, *J* is longer than *I* by exactly two 1's. If there is no upper bound on the length of *I*, this sort of restriction will give rise to a non-regular set (this follows from a theorem of Nerode; see Rabin and Scott 1959 and Chomsky and Miller 1958).

Phonological rules do, then, occasionally exceed the capacity of finite-state machines. Apparently, however, this happens only in the restricted sort of case just considered, where integral coefficients are being referred to or manipulated.

(76c) has been given as right-linear but will work equally well when applied iteratively. The rule given by Schachter and Fromkin (p. 108) is in fact meant to be iterative and is essentially the same

as our rule, though differing in its outward form and corresponding rather to (76) as a whole. Notice that there is no way at all to express the rule under the simultaneous mode of application. For suppose we attempted a simultaneous formulation. The determination of the pitch of each syllable would have to be made independently of the other syllabics, since the context of each syllabic must be referred to in original input form. What we would have to say, then, is this:

- (i) A +tone syllabic receives *n* pitch when preceded in the phrase by exactly *n*-1 occurrences of sequences subsumed under [+tone] [-syl]* [+syl, -tone].
- (ii) A -tone syllabic receives *n*+3 pitch when preceded in the phrase by exactly *n* occurrences of sequences subsumed under [+syl, -tone] [+syl]* [+tone].

However, the schematic notation as it now stands is unable to express a dependency between the number of occurrences of a particular kind of phonological string and the size of an integer coefficient, at least when there is no upper bound on these values. The most it can do in this respect is to state dependence between integers. We could, therefore, formulate a simultaneous version of (76c) only if we introduced some further device, such as the notation *X*^{*n*} meaning a string of exactly *n* *X*'s. It would then be imperative to allow *I* to be a variable as well as a specific integer. However, right-linear application obviates the necessity for this extra device in the downdrift rule.

There is one other notable example of a feature that has been said to take on any positive integer as a coefficient. This feature is stress. In SPE the integer coefficients designate degrees of strength, 1 denoting the strongest stress, 2 the next strongest stress, and so on. Stress has some peculiar properties which set it apart from other features and appear to call for special conventions.

In many cases the introduction of a 1 stress causes weakening of all previously present integrally valued stresses by one degree. Thus the English alternating stress rule (rule 17, Chapter 5, in SPE) will not only place a 1 stress on the antepenultimate vowel

of a final-stressed word but will also weaken the previously present final 1 stress to 2 stress. Thus *cavalcāde* becomes *cāvalcāde*. Technically, we could treat the weakening of the final stress as a separate rule. However, weakening of this sort is such a frequent concomitant of primary stress placement that the two phenomena seem to be part of one process. Consequently Chomsky and Halle have decided in SPE to let this weakening take place automatically. If we incorporate their convention about stress into our formalism we can write the alternating stress rule simply as

$$\begin{bmatrix} +sy1 \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} 1 \text{ stress} \\ (\sigma) \end{bmatrix} / \$* - [-sy1]^* [+sy1] [-sy1]^* [1 \text{ stress}] [-sy1]^*$$

Many rules that introduce a primary stress operate according to a principle of disjunctive application. Consider a revised version of the alternating stress rule proposed by Ross (1969). For simplicity of illustration we take Ross' preliminary revision on page 9, omitting boundary symbols. Ross' rule can be thought of as having the two ordered subcases (77a) and (77b) (some slight modifications have been made to fit the rule to our notation).

(77)

$$\begin{aligned} \text{(a)} \quad & \begin{bmatrix} +sy1 \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} 1 \text{ stress} \\ (\sigma) \end{bmatrix} / \$* - [-sy1]^* [+sy1] [-sy1]^* [1 \text{ stress}] [-sy1] \\ \text{(b)} \quad & \begin{bmatrix} +sy1 \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} 1 \text{ stress} \\ (\sigma) \end{bmatrix} / \$* - [-sy1]^* [1 \text{ stress}] [-sy1]^* \end{aligned}$$

In general, all nouns are subject to (77) (most verbs are exempted from (77b) by a redundancy rule in Ross' view). However, the two cases of (77) are applied disjunctively: if a noun receives a 1 stress by (77a) it is no longer subject to (77b). Thus from *cavalcāde* and *kāvāde*, where the final 1 stress has been introduced previously by the main stress rule, we obtain *cāvalcāde* and *kāvāde*. It is the principle of disjunctive application that prevents the derivation

$$\begin{array}{ll} \text{cavalcāde} & \\ \text{cāvalcāde} & (77a) \\ *cāvālcāde & (77b) \end{array}$$

It is frequently the case that the second of two disjunctively ordered rules omits material present in the first rule. This is true in (77), for example. An equivalent of (77b) can be obtained from (77a) by simply omitting the [+sy1] from the expression to the right of the slash. This fact has led to the association of disjunctive ordering with parentheses and with the closely allied angle bracket notation. Thus (77) would be given in abbreviated form as (78).

(78)

$$\begin{bmatrix} +sy1 \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} 1 \text{ stress} \\ (\sigma) \end{bmatrix} / \$* - [-sy1]^* ([-sy1]) [-sy1]^* [1 \text{ stress}] [-sy1]^*$$

Although parenthesis notation gives the correct result in many cases there are certain reasons why we will not incorporate it officially into our formalism. For one thing parentheses are not formally associated with the stress feature. Indeed, with their aid we can formulate many disjunctively applied rules that have nothing to do with stress. In this respect the principle of disjunctive ordering is grossly over-generalized, for it seems to be properly associated only with certain rules which introduce a primary accent (whether of pitch or stress). Supposed cases of disjunctive ordering that do not fall under this rubric appear to be spurious. A few illustrative examples will suggest why I believe this to be so.¹

¹ Chomsky and Halle (p. 366 of SPE) are aware of the sort of problem about to be discussed and offer some preliminary suggestions for solving it. We have not followed up these suggestions because the problem does not arise within the rule formalism we are developing here. It must be granted that the case discussed by Chomsky and Halle, a rule of Latvian, presents a difficulty of another kind to our formalism. Adapted slightly to our notation, but retaining the parentheses and morpheme boundary of Chomsky and Halle's formulation, the rule would appear as follows:

$$\begin{bmatrix} +sy1 \\ +high \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} -sy1 \\ (\sigma) \end{bmatrix} / \$* - (+) [+sy1] \$*$$

This is taken to represent the following sequence of rules:

$$\text{(a)} \quad \begin{bmatrix} +sy1 \\ +high \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} -sy1 \\ (\sigma) \end{bmatrix} / \$* - + [+sy1] \$*$$

One of the rules proposed in SPE (23III in Chapter 5) replaces lax *u* with tense *i* before a string subsumed under $[-\text{syll}]-\text{cons}]$ or $[-\text{cons}]$. This appears to cry out for parenthesis notation, and the authors of SPE do in fact use an equivalent notation involving a subscript 0 and a superscript 1. Adapted slightly to fit our formalism more closely, the rule could be given as (79).

$$(79) \begin{bmatrix} +\text{syll} \\ -\text{tns} \\ +\text{back} \\ +\text{high} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} +\text{tns} \\ -\text{rnd} \\ (\sigma) \end{bmatrix} / \text{S}^* - [-\text{syll}]_0^1 [-\text{cons}]\text{S}^*$$

According to the conventions in SPE (pp. 61-2, 199), (79) would be expanded into the two ordered rules of (80),

- (80) (a) $u \rightarrow i / \text{S}^* - [-\text{syll}]-\text{cons}]\text{S}^*$
 (b) $u \rightarrow i / \text{S}^* - [-\text{cons}]\text{S}^*$

Now consider the effect of the $u \rightarrow i$ rule on the word *izuel*. This word presumably has the underlying phonological representation *uzuel* (SPE, p. 228). S-voicing apparently should apply before the $u \rightarrow i$ rule since s-voicing precedes velar softening according to SPE, p. 221, and velar softening precedes the $u \rightarrow i$ rule 23III of Chapter 5. The input to the $u \rightarrow i$ rule will, then, be *izuel*. This string has two vowels which fit the conditions of the $u \rightarrow i$ rule, the first *u* satisfying case (80a) and the second *u* satisfying case (80b). According to the conventions of SPE, however, the two cases are applied disjunctively; consequently only the first case (80a) will apply, and the output will be *izuel*. Unfortunately,

$$(b) \begin{bmatrix} +\text{syll} \\ +\text{high} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} -\text{syll} \\ (\sigma) \end{bmatrix} / \text{S}^* - [+ \text{syll}]\text{S}^*$$

When applied in the indicated order (a) and (b) will correctly convert *kur^w+tal* and *ait^w+a* into *kur^w+yai* and *ait^w+a*, respectively. As our formalism now stands we cannot, as Chomsky and Halle can, regard (a) and (b) as subcases of a single rule. If they are subcases of a single right-linear or simultaneous rule they yield **ait^w+a* instead of *ait^w+a*; if they are subcases of a single left-linear rule, they yield **kur^w+yai* instead of *kur^w+yai*.

though, this output cannot yield the correct pronunciation of the word, for the remainder of the derivation would be as follows (numbers not in parentheses refer to rules in Chapter 5 of SPE):

| | |
|---------------------|------|
| <i>i</i> zuel | 23IV |
| <i>i</i> zuel | 29 |
| <i>y</i> izuel | 31 |
| <i>y</i> iwzūwael | 34 |
| * <i>y</i> ūwzūwael | 43 |

The desired phonetic form is *yūwzūwael*, with palatalization of the *z*. But it is just this correct form that will be obtained of the two cases (80a) and (80b) of the $u \rightarrow i$ rule are applied conjunctively or simultaneously. Then the input string *izuel* will be converted to *izuel* and the remainder of the derivation will proceed as desired:

| | |
|--------------------|----|
| <i>i</i> zuel | 29 |
| <i>y</i> izuel | 31 |
| <i>y</i> iwzyūwael | 34 |
| <i>y</i> ūwzyūwael | 37 |
| <i>y</i> ūwzūwael | 38 |
| <i>y</i> ūwzūwael | 43 |

Note that within our formalism, where there is no explicit analogue to the notion of ordered expansion of rules, (79) will be interpreted correctly if the notation $[-\text{syll}]_0^1$ is regarded as equivalent to $\{[-\text{syll}], \emptyset\}$ and if (79) is regarded as a simultaneous, right-linear, or left-linear rule. For then (79) will convert *izuel* to *izuel* by virtue of the application

$$\emptyset \begin{matrix} \dot{u} & z & u \\ \dot{i} & z & i \end{matrix} \text{ael}$$

Variable notation too has been interpreted according to the principle of disjunctive ordering, with equally disastrous con-

sequences. Consider the second part of the vowel shift (rule 33 of Chapter 5 in SPE), which could be given as

$$(81) \begin{bmatrix} X \\ \text{flow} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} -\text{flow} \\ (\sigma) \end{bmatrix} / \$* - \$*,$$

where X is the expression

+syl
+tns
+stress
-high
yback
yround

According to the convention described on p. 357 of SPE, (81) must be regarded as standing for the disjunctive sequence (82).

$$(82) \quad (a) \begin{bmatrix} X \\ +\text{low} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} -\text{low} \\ (\sigma) \end{bmatrix} / \$* - \$*$$

$$(b) \begin{bmatrix} X \\ -\text{low} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} +\text{low} \\ (\sigma) \end{bmatrix} / \$* - \$*$$

Now consider the effect on the word *migrate*. The underlying representation is presumably *migræt* (in SPE, p. 144, this word is thought to have a + boundary before the *x*, but this fact is of no relevance here except for its effect on the stress). The stress rules, diphthongization and the first part of the vowel shift rule will convert *migræt* into the intermediate form *mêygrêyt*. Applying case (82a) of the second part of the vowel shift we obtain *mêygrêyt*. Now, however, because of disjunctive ordering, (82b) cannot now apply to this output. But this is not what is wanted; the desired output is *mêygrêyt*. Conjunctive ordering of (82a) and (82b) will not do either, for then the output *mêygrêyt* will be obtained. Within the framework of SPE the most appropriate method of application seems to be the simultaneous; for then (82a) will apply

to the second vowel but not to the first, and (82b) will apply to the first vowel but not to the second, as required. Within our formalism the correct result is obtained immediately if (81) is regarded as simultaneous, right-linear, or left-linear. Then the rule will convert *mêygrêyt* into *mêygrêyt* by virtue of the application

$$\begin{matrix} \acute{e} & \grave{e} \\ m & \grave{a} & ygr & \acute{e} \\ & \grave{e} & & t \end{matrix}$$

There are several more cases in SPE where the notation incorrectly calls for disjunctive ordering; the amusing game of finding them is left to the reader. There are, in addition, a number of cases where disjunctive application is called for by the notation but is neither supported nor invalidated by any crucial evidence in SPE. Rule 24 of Chapter 5 in SPE is a particularly interesting example. The first part of the rule would take the following form within our notation:

$$(83) \quad Y(Z)X \begin{Bmatrix} 0 \\ W \end{Bmatrix} : K$$

where

$$Y = \begin{bmatrix} \alpha\text{stress} \\ +\text{syl} \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} 11 \text{ stress} \\ (\sigma) \end{bmatrix} / [-\text{stress}]^* - [-\text{syl}]^*,$$

$$Z = \begin{bmatrix} +\text{syl} \\ -\text{tns} \\ -\text{stress} \end{bmatrix} ([-\text{syl}]^*),$$

$$X = [-\text{syl}]^* \begin{bmatrix} +\text{syl} \\ \beta\text{stress} \end{bmatrix} [-\text{syl}]^* \begin{bmatrix} +\text{syl} \\ \gamma\text{stress} \end{bmatrix} [-\text{syl}]^*,$$

$$W = \begin{bmatrix} \delta\text{stress} \\ +\text{syl} \end{bmatrix} \$*,$$

and K is the condition

($-(\alpha = 1)$ and $-(\beta = 1)$ and $-(\beta = 11)$ and $-(\gamma = 1\delta)$)
(that is, α is not 1, β is weaker than 2, and δ is weaker than γ).

all integrally valued stresses in Ω , including the 0 stresses, are weakened by one.

With these conventions we can now formulate Ross' version of the alternating stress rule as follows:

$$(86) \quad R: \begin{bmatrix} +syI \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} 0 \text{ stress}^s \\ (\sigma) \end{bmatrix} / \$* - [-syI]^* \left\{ \begin{array}{l} [+syI] \\ [-syI]^* \end{array} \right. \\ \left. \left[1 \text{ stress} \right] [-syI]^* \right.$$

Application of (86) to $kavək\bar{a}d$ (perhaps the representation of *cavalcade* at this stage of derivation) will take the following form:

$$\begin{array}{c} \bar{a} \\ k \quad 0 \\ \bar{a} \end{array} \begin{array}{c} 1 \\ vək\bar{a}d \\ \bar{a} \end{array}$$

The output $k\bar{a}vək\bar{a}d$ is turned into $k\bar{a}vək\bar{a}d$ by convention (85c). Notice that (85) cannot give the output $k\bar{a}vək\bar{a}d$ by virtue of any putative application

$$\begin{array}{c} \bar{a} \quad \bar{a} \\ k \quad v \quad 0 \\ \bar{a} \quad \bar{a} \end{array} \begin{array}{c} 1 \\ k\bar{a}d \\ \bar{a} \end{array}$$

The reason is that the supposed fourth step,

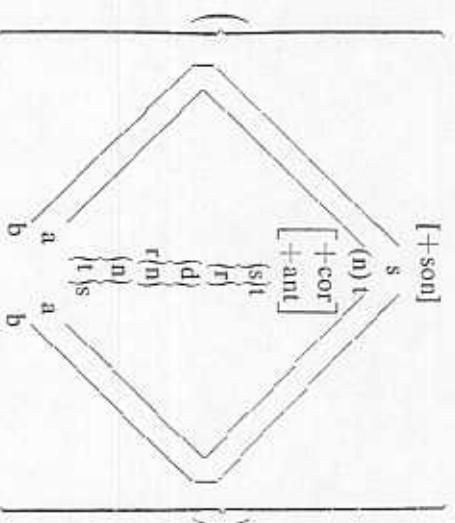
$$\bar{a} \rightarrow \bar{a}^0/k\bar{a}v\bar{a}^0 - 1k\bar{a}d,$$

though subsumed under the schema in (86), is not a subrule of (85), being excluded by convention (85b).

Ross' version of the main stress rule of English, though not the version in SPE, fits quite well into the framework proposed here. Adapted slightly to our notation, Ross' rule as given on p. 45 takes the form of (87).

$$(87) \quad \begin{bmatrix} +syI \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} 1 \text{ stress} \\ (\sigma) \end{bmatrix} / \$* - [-syI]^* \left(\begin{bmatrix} +syI \\ -ins \end{bmatrix} \left([-syI] \right) \begin{bmatrix} +syI \\ -ins \end{bmatrix} \right) \left[K \right] L$$

where K is the expression



and L is the right syntactic bracket

$$\left. \right] b \langle a \langle N \rangle_a A \rangle b$$

According to the conventions of SPE (cf. especially the Appendix to Chapter 8), rule (87) will have the partial expansion shown in (88).

$$(88) \quad \begin{array}{l} \text{a) } \begin{bmatrix} +syI \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} 1 \text{ stress} \\ (\sigma) \end{bmatrix} / \$* - [-syI]^* \begin{bmatrix} +syI \\ -ins \end{bmatrix} \left([-syI] \right) \begin{bmatrix} +syI \\ -ins \end{bmatrix} \left[K \right] L \\ \text{b) } \begin{bmatrix} +syI \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} 1 \text{ stress}^s \\ (\sigma) \end{bmatrix} / \$* - [-syI]^* \begin{bmatrix} +syI \\ -ins \end{bmatrix} \left[K \right] L \\ \text{c) } \begin{bmatrix} +syI \\ \sigma \end{bmatrix} \rightarrow \begin{bmatrix} 1 \text{ stress} \\ (\sigma) \end{bmatrix} / \$* - [-syI]^* L \end{array}$$

The three cases (88a), (88b), and (88c) of rule (87) stress antepenultimate, penultimate, and final vowels, respectively. Since the cases are applied disjunctively in the order given, the effect is to stress the leftmost vowel that fits any of the three cases. We can, then, adapt (87) completely to our formalism with very minor

revisions, provided we adopt the customary angle bracket notation and the usual way of referring to syntactic categories. All we need to do is to make (87) a right-linear rule, substitute the expression [0 stress, (σ)] for [i stress, (σ)] to the right of the arrow, and replace each expression of the form (U) with {U, \emptyset }.

There are a few cases where disjunctive application seems appropriate and is in fact achievable under our conventions but not under those of SPE. According to Harris (1968: 74), Komi Jazva has a word-level rule that places stress on the rightmost vowel that is not preceded anywhere within the word by a tense vowel. Thus if there is a tense vowel, the first tense vowel receives stress, while if there is no tense vowel, then the last vowel receives stress. With our conventions we would write the rule as follows:

$$(89) \quad L: \left[\begin{array}{c} +syl \\ \sigma \end{array} \right] \rightarrow \left[\begin{array}{c} 0 \text{ stress} \\ (\sigma) \end{array} \right] / \left[\left\{ \begin{array}{c} -syl \\ -ins \end{array} \right\} \right] * - \$ *$$

There is no way in the system of SPE to achieve the effect of disjunctive application here, since in SPE the star notation is associated with simultaneous application. Parentheses around the left-hand starred expression will not do either, since the first of the two disjunctively ordered cases will be (89) itself.

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