Enlargement of Calderbank Shor Steane quantum codes

Andrew M. Steane*

February 1998

Abstract

It is shown that a classical error correcting code $C = [n, k, d]$ which contains its dual, $C^{\perp} \subseteq C$, and which can be enlarged to $C' = [n, k' > k+1, d']$, can be converted into a quantum code of parameters $[[n, k+k'-n, \min(d, \lceil 3d'/2 \rceil)]]$. This is a generalisation of a previous construction, it enables many new codes of good efficiency to be discovered. Examples based on classical Bose Chaudhuri Hocquenghem (BCH) codes are discussed.

keywords Quantum error correction, BCH code, CSS code

*Department of Physics, University of Oxford, Clarendon Laboratory, Parks Road, Oxford OX1 3PU, England.

Quantum information theory is rapidly becoming a well-established discipline. It shares many of the concepts of classical information theory but involves new subtleties arising from the nature of quantum mechanics [1]. Among the central concepts in common between classical and quantum information is that of error correction, and the error correcting code. Quantum error correcting codes have progressed from their initial discovery [2, 3] and the first general descriptions [4, 3, 5] to broader analyses of the physical principles [8, 10, 11, 13] and various code constructions [11, 13, 17, 9, 12, 16, 14, 15]. A thorough discussion of the principles of quantum coding theory is offered in [6], and many example codes are given, together with a tabulation of codes and bounds on the minimum distance for codeword length n up to $n = 30$ quantum bits.

For larger n there is less progress, and only a few general code constructions are known. The first important quantum code construction is that of [4, 3, 5], the resulting codes are commonly referred to as Calderbank Shor Steane (CSS) codes. It can be shown that efficient CSS codes exist as $n \to \infty$, but on the other hand these codes are not the most efficient possible. I will present here a method which permits most CSS codes to be enlarged, without an attendant reduction in the minimum distance of the code. The resulting codes are therefore more efficient than CSS codes. The examples I will give are found to be among the most efficient quantum codes known, and enable some of the bounds in [6] to be tightened. The code construction is essentially the same as that described for Reed-Muller codes in [16], the new feature is to understand how the method works and thus prove that it remains successful for a much wider class of code. After this some relevent theory of Bose Chaudhuri Hocquenghem (BCH) codes [18, 19, 21] will be given and used to construct a table of example quantum codes built by the new method. The codes are *additive* and *pure* in the nomenclature of [6]. A pure additive code is *nondegenerate* in the nomenclature of [13].

1 Quantum coding

Following [6], the notation $[[n, k, d]]$ is used to refer to a quantum error correcting code for n qubits having 2^k codewords and minimum distance d. Such a code enables the quantum information to be restored after any set of up to $\lfloor (d-1)/2 \rfloor$ qubits has undergone errors. In addition, when d is even, $d/2$ errors can be detected. We restrict attention to the 'worst case' that any defecting qubit (ie any qubit undergoing an unknown interaction) might change state in a completely unknown way, so all the error processes X, Z and $Y = XZ$ must be correctable [5, 7, 8, 13].

A quantum error correcting code is an eigenspace of a commutative subgroup of the group E of tensor products of Pauli matrices. The commutativity condition can be expressed [11, 6, 16]

$$H_x \cdot H_z^T + H_z \cdot H_x^T = \mathbf{0}. \tag{1}$$

where H_x and H_z are $(n-k \times n)$ binary matrices which together form the *stabilizer* $\mathcal{H} = (H_x | H_z)$. All vectors $(u_x|u_z)$ in the code (where u_x and u_z are n-bit strings) satisfy $H_x \cdot u_z + H_z \cdot u_x = 0$. These are generated by the generator $\mathcal{G} = (G_x | G_z)$ which therefore must satisfy

$$H_x \cdot G_z^T + H_z \cdot G_x^T = \mathbf{0}. \tag{2}$$

In other words $\mathcal{H}$ may be obtained from $\mathcal{G}$ by swapping the X and Z parts, and extracting the dual of the resulting $(n+k) \times 2n$ binary matrix. The rows of G_x and G_z have length n, and the number of rows is $n+k$.

The weight of a vector $(u_x|u_z)$ is the Hamming weight of the bitwise OR of u_x with u_z. The minimum distance d of the code $\mathcal{C}$ is the largest weight such that there are no vectors of weight $< d$ in $\mathcal{C} \setminus \mathcal{C}^\perp$, where the dual is with respect to the inner product $((u_x|u_z),(v_x|v_z)) \equiv u_x \cdot v_z + u_z \cdot v_x$. A *pure* code has furthermore no vectors of weight $< d$ in $\mathcal{C}$, apart from the zero vector.

The CSS code construction [4, 5] is to take classical codes C_1 and C_2 with $C_1^\perp \subseteq C_2$, and form

$$\mathcal{G} = \left(\begin{array}{c|c} G_1 & 0 \\ 0 & G_2 \end{array} \right), \qquad \mathcal{H} = \left(\begin{array}{c|c} H_2 & 0 \\ 0 & H_1 \end{array} \right). \tag{3}$$

where G_i and H_i are the classical generator and check matrices. The dual condition $C_1^\perp \subseteq C_2$ ensures that $H_1 \cdot H_2^T = H_2 \cdot H_1^T = 0$ and therefore the commutativity condition (1) is satisfied. If $C_1 = [n, k_1, d_1]$ and $C_2 = [n, k_2, d_2]$ then the minimum distance of the quantum code is $\min(d_1, d_2)$ and the number of rows in $\mathcal{G}$ is $k_1 + k_2$, leading to quantum code parameters $[[n, k_1 + k_2 - n, \min(d_1, d_2)]]$.

An interesting subset of CSS codes is that given by the above construction starting from a classical $[n, k, d]$ which contains its dual, leading to a quantum $[[n, 2k-n, d]]$ code.

2 New code construction

I will present the new construction by stating and proving the following.

Theorem 1. *Given a classical binary error correcting code $C = [n, k, d]$ which contains its dual, $C^\perp \subseteq C$, and which can be enlarged to $C' = [n, k' > k+1, d']$, a pure quantum code of parameters $[[n, k+k'-n, \min(d, \lceil 3d'/2 \rceil)]]$ can be constructed.*

Proof. The generator for the quantum code is

$$\mathcal{G} = \left(\begin{array}{c|c} D & AD \\ G & 0 \\ 0 & G \end{array} \right), \tag{4}$$

where G generates the classical code C, and G and D together generate C', as does G and AD together (we will choose A such that D and AD generate the same set). From the dual conditions specified in the theorem, $G \cdot G^T = 0$, $G \cdot (AD)^T = 0$, $D \cdot G^T = 0$ and $D \cdot (AD)^T = 0$, therefore $\mathcal{G}_x \cdot \mathcal{G}_z = 0$, from which $\mathcal{H}_x \cdot \mathcal{H}_z = 0$ and the commutativity condition (1) is satisified.

Since the number of rows in the generator is $k+k'$, the dimension of the quantum code is $k+k'-n$. It remains to prove that the minimum distance is $\min(d, \lceil 3d'/2 \rceil)$.

We choose A such that D and AD generate the same set. Therefore for any vector $(u|v)$ generated by $(D|AD)$, either $u = v$ or $\mathrm{wt}(u+v) \geq d'$. We choose the map A such that $u = v$ never occurs (a fixed point free map). This can be achieved as long as D has more than one row, by, for example, the map

$$A = \begin{pmatrix} 0100\ldots0 \\ 0010\ldots0 \\ 0001\ldots0 \\ \ldots \\ 0000\ldots1 \\ 1100\ldots0 \end{pmatrix}. \tag{5}$$

To complete the proof we will show that for any non-zero vector $(u|v)$ generated by $\mathcal{G}$, $\mathrm{wt}(u|v) \geq \min(d, 3d'/2)$ (and therefore $\mathrm{wt}(u|v) \geq \min(d, \lceil 3d'/2 \rceil)$.)

For the non-zero vector $(u|v)$, if either $\mathrm{wt}(u) \geq d$ or $\mathrm{wt}(v) \geq d$ then $\mathrm{wt}(u|v) \geq d$, so the conditions of the theorem are satisfied. The only remaining vectors are those for which both $\mathrm{wt}(u) < d$ and $\mathrm{wt}(v) < d$. Now, $\mathrm{wt}(u)$ can only be less than d if D is involved in the generation of u, and $\mathrm{wt}(v)$ can only be less than d if AD is involved in the generation of v, since G on its own generates a binary code of minimum distance d. However, since the map A is fixed-point free, and using the fact that D and AD generate the same set, the binary vector $u + v$ is not zero and is a member of a distance d' code, therefore $\mathrm{wt}(u+v) \geq d'$. We thus have the conditions $\{\mathrm{wt}(u) \geq d'$, $\mathrm{wt}(v) \geq d'$, $\mathrm{wt}(u+v) \geq d'\}$. These are sufficient to imply that $\mathrm{wt}(u|v) \geq 3d'/2$. For, if u and v overlap in p places, then $\mathrm{wt}(u+v) = \mathrm{wt}(u) - p + \mathrm{wt}(v) - p$ and $\mathrm{wt}(u|v) = \mathrm{wt}(u) + \mathrm{wt}(v) - p = (\mathrm{wt}(u) + \mathrm{wt}(v) + \mathrm{wt}(u+v))/2 \geq 3d'/2$. This completes the proof.

The above construction was applied to Reed-Muller codes in [16]. These codes are not very efficient (they have small k/n for given n, d) but they have the advantage of being easily decoded. A large group of classical codes which combine good efficiency with ease of decoding are the BCH codes. They include Reed Solomon codes as a subset. I will now derive a set of quantum error correcting codes from binary BCH codes using the above construction, combined with some simple BCH coding theory.

3 Application to binary BCH codes

Properties of BCH codes are discussed and proved in, for example, [21]. A binary BCH code of designed distance δ is a cyclic code of length n over GF(2) with generator polynomial

$$g(x) = \mathrm{l.c.m.}\{M^{(b)}(x), M^{(b+1)}(x), \ldots, M^{(b+\delta-2)}(x)\} \tag{6}$$

where

$$M^{(s)}(x) = \prod_{i \in C_s} \left(x - \alpha^i\right), \tag{7}$$

in which α is a primitive nth root of unity over GF(2), and C_s is a cyclotomic coset mod n over GF(2), defined by

$$C_s = \{s, 2s, 4s, \ldots, 2^{m_s - 1}s\}, \tag{8}$$

where $m_s = |C_s|$ is obtained from $2^{m_s}s \equiv s \bmod n$. The dimension of the code is $k = n - \deg(g(x))$. From (6) and (7) this implies $k = n - \sum_s |C_s|$ where the sum ranges from $s = b$ to $s = \delta + b - 2$ but only includes each distinct cyclotomic coset once. This can also be expressed $k = n - |\mathcal{I}_C|$ where $\mathcal{I}_C = C_b \cup C_{b+1} \cup \cdots \cup C_{b+\delta-2}$ is called the *defining set*. The minimum distance of the code is $d \geq \delta$.

The dual of a cyclic code is cyclic. Grassl *et al.* [20] derive the useful criterion that a cyclic code contains its dual if the union of cyclotomic cosets contributing to $g(x)$ does not contain both C_s and C_{n-s}. In other words

$$\{(n-i) \notin \mathcal{I}_C \ \ \forall i \in \mathcal{I}_C\} \Rightarrow C^{\perp} \subseteq C. \tag{9}$$

3.1 Primitive BCH codes

Consider first the BCH codes with $n = 2^m - 1$, the so-called primitive BCH codes. In order to find the codes which satisfy the condition (9), we will find the smallest s such that $n - r \in C_s$ for some $r \leq s$. The largest permissible designed distance will then be $\delta = s$. For even m, the choice $s = 2^{m/2} - 1$ gives $s2^{m/2} = n - s \Rightarrow C_s = C_{-s}$, so this is an upper bound on s. For odd m, an upper bound is provided by $s = 2^{(m+1)/2} - 1$ since then $s2^{(m-1)/2} = n - (s-1)/2$. We will show that these upper bounds can be filled, i.e. that no smaller s leads to $n - r \in C_s$ for $r \leq s$.

For $n = 2^m - 1$, the elements of the cyclotomic cosets C_s are largest when s is one less than a power of 2, $s = 2^j - 1$. Specifically, for $s = 2^j - 1$ we have $(s2^i \bmod n) > (r2^i \bmod n)\ \forall i < m, r < s$. This is obvious for $s2^i < n$ and the proof for $s2^i > n$ is straightforward. The largest element in C_s $(s = 2^j - 1)$ is obtained for the largest i such that $s2^i < n$, giving $\max(C_s) = 2^m - 2^{m-j} = n - r$ where $r = 2^{m-j} - 1$. This element $\max(C_s) = n - r$ is the largest in the defining set $\mathcal{I}_C$ for a code of designed distance $\delta = s$, therefore it is only possible for $\mathcal{I}_C$ to contain both i and $n - i$ (for any i) if it contains r and $n - r$, since $r = 2^{m-j} - 1$ is the smallest element in its coset, and any other pairs $i, n - i$ must have $i > r$. Finally, we have a failure of the condition (9) only if $r \leq s$, that is $2^{m-j} - 1 \leq 2^j - 1$, therefore $j \geq \lceil m/2 \rceil$.

To summarise the above, we have the proved following:

Lemma: *The primitive binary BCH codes contain their duals if and only if the designed distance satisfies*

$$\delta \leq 2^{\lceil m/2 \rceil} - 1 \tag{10}$$

Using the code construction of theorem 1, together with this lemma, the list of quantum codes in table 1 is obtained. The further property used is that BCH codes are nested, i.e. codes of smaller distance contain those of larger, which is obvious since the former can be obtained from the latter by deleting parity checks. The first entry for each value of n uses $\{C =$ extended BCH code$\}$ with $\{C' =$ even weight code$\}$ to obtain a distance 3 quantum code. The codes of larger distance involve only BCH codes, for these a quantum code is obtained both from the unextended and extended versions. The parameters $[[n, K, D]]$ given in the table are for the extended BCH codes (i.e. extended by an overall parity check). Using unextended codes leads to a further quantum code of parameters $[[n-1, K+1, D-1]]$, for $D > 3$.

3.2 Non-primitive BCH codes

When $n \neq 2^m - 1$ the cyclotomic cosets mod n do not have so much structure so in general the only way to find if condition (9) is satisfied is to examine each coset individually.

One way in which the requirement (9) is not met is if C_s contains both i and $-i$ mod n, which implies $C_s = C_{-s}$, for some $C_s \subseteq \mathcal{I}_C$. If s is the smallest element in C_s, then $i, n-i \in C_s$ if and only $s, n-s \in C_s$, from which $s2^j \equiv -s \bmod n$ for some $j < m_s$. Multiplying by 2^j we have $s2^{2j} \equiv -s2^j \equiv s \bmod n$, therefore $j = m_s/2$ and this is only possible for even m_s. Furthermore, since $m_{s \geq 1}$ is a factor of m_1, m_s can be even only if m_1 is even. This observation slightly reduces the amount of checking to be done.

The values of n in the range $1 < n \leq 127$ for which C_1 does not contain $n-1$ are { 7,15,21,23,31,35,39,45, 47,49,51,55,63,69,71,73,75,77,79,85,87,89,91,93,95,103, 105,111,115,117,119,121,123,127}. An efficient code is obtained if one or more of the cosets is small, this happens for $n =$ 21, 23, 45, 51, 73, 85, 89, 93, 105, 117 (not counting primitive codes). Quantum codes obtained from BCH codes with these values of n are listed in table 2. Further good codes exist in the range $127 < n < 511$ for $n =$ 133, 151, 153, 155, 165, 189, 195, 217, 219, 255, 267, 273, 275, 279, 315, 337, 341, 381, 399, 455.

4 Efficiency

The code parameters in tables 1 and 2 compare well with the most efficient quantum codes known. For example, the $[[22, 5, 6]]$ code permits some of the lower existence bounds in [6] to be raised, and the $[[32, 15, 6]]$ and $[[32, 5, 8]]$ codes fill lower existence bounds. The $[[93,68,5]]$ code is comparable with the $[[85,61,5]]$ code quoted in [6], though the $[[93,53,7]]$ code is not as good as $[[85,53,7]]$ quoted in [6]. Obviously the quantum codes based on BCH codes will be best for primitive BCH codes, so we expect the codes in table 1 rather than table 2 to compare best with other code constructions. Indeed, the distance 3 codes in table 1 are the previously known Hamming codes [13, 16, 6] and are optimal.

The quantum codes constructed by theorem 1 have an upper bound on the rate $K/n = (k+k')/n - 1$ arising from the upper bound on k and k' for binary codes. In the asymptotic limit this bound on the quantum codes is

$$K/n < R(d/n) + R(2d/3n) - 1, \tag{11}$$

where $R(d/n)$ is the maximum rate of a binary $[n,k,d]$ linear code. For example the sphere-packing bound is $R(d/n) < 1 - H(d/2n)$; the codes we have discussed have parameters lying close to this bound (though in the limit of large n it is known that BCH codes are no longer efficient). Taking $R(x)$ equal to the McEliece-Rodemich-Rumsey-Welch upper bound [22], we find $K/n \geq 0$ for $d/n < 0.2197$ in the limit of large n. This may be compared with $d/n < 0.1825$ for CSS codes and the limit $d/n < 0.308$ for pure quantum stabilizer codes discussed by Ashikhmin [23].

The author is supported by the Royal Society and by St Edmund Hall, Oxford.

References

[1] For a general introduction to quantum information theory, see, for example, A. M. Steane, "Quantum computing", *Rep. Prog. Phys.*, to be published (preprint quant-ph/9708022).

[2] P. W. Shor, "Scheme for reducing decoherence in quantum computer memory," *Phys. Rev. A*, vol. 52, pp. R2493-R2496, Oct. 1995.

[3] A. M. Steane, "Error correcting codes in quantum theory," *Phys. Rev. Lett.*, vol. 77, pp. 793-767, July 1996.

[4] A. R. Calderbank and P. W. Shor, "Good quantum error-correcting codes exist," *Phys. Rev. A*, vol. 54, pp. 1098-1105, Aug. 1996.

[5] A. M. Steane, "Multiple particle interference and quantum error correction," *Proc. Roy. Soc. Lond. A*, vol. 452, pp. 2551-2577, Nov. 1996.

[6] A. R. Calderbank, E. M. Rains, P. W. Shor and N. J. A. Sloane, "Quantum error correction via codes over $GF(4)$", *IEEE Trans. Information Theory*, to be published. (LANL eprint quant-ph/9608006).

[7] A. Ekert and C. Macchiavello, "Quantum error correction for communication," *Phys. Rev. Lett.*, vol. 77, pp. 2585-2588, Sept. 1996.

[8] E. Knill and R. Laflamme, "A theory of quantum error correcting codes," *Phys. Rev. A*, vol. 55, pp. 900-911, (1997).

[9] R. Laflamme, C. Miquel, J. P. Paz and W. H. Zurek, "Perfect quantum error correcting code," *Phys. Rev. Lett.*, vol. 77, pp. 198-201, July 1996.

[10] C. H. Bennett, D. P. DiVincenzo, J. A. Smolin and W. K. Wootters, "Mixed state entanglement and quantum error correction," *Phys. Rev. A*, vol. 54, pp. 3822-3851, (1996).

[11] A. R. Calderbank, E. M. Rains, N. J. A. Sloane and P. W. Shor, "Quantum error correction and orthogonal geometry," *Phys. Rev. Lett.*, vol. 78, pp. 405-409, (1997).

[12] A. M. Steane, "Simple quantum error correcting codes," *Phys. Rev. A*, vol. 54, pp. 4741-4751, (1996).

[13] D. Gottesman, "Class of quantum error-correcting codes saturating the quantum Hamming bound," *Phys. Rev. A*, vol. 54, pp. 1862-1868, (1996).

[14] E. M. Rains, "Nonbinary quantum codes," Preprint quant-ph/9703048.

[15] E. M. Rains, "Quantum codes of minimum distance two," Preprint quant-ph/9704043.

[16] A. M. Steane, "Quantum Reed-Muller codes," *IEEE Trans. Inf. Theory*, to be published (preprint quant-ph/9608026).

[17] D. Gottesman, "Pasting quantum codes," Preprint quant-ph/9607027.

[18] R. C. Bose and C. R. Ray-Chaudhuri, "On a class of error-correcting binary group codes," *Information and Control*, vol. 3, pp/ 68-79.

[19] A. Hocquenghem, "Codes correcteurs d'erreurs," *Chiffres*, vol 2, pp. 147-156, September, 1959.

[20] M. Grassl, Th. Beth and T. Pellizzari, "Codes for the quantum erasure channel," *Phys. Rev. A*, vol 56, pp. 33-38 (1997).

[21] F. J. MacWilliams and N. J. A. Sloane, *The Theory of Error-Correcting Codes,* Amsterdam: North-Holland, 1977.

[22] R. J. McEliece, E. R. Rodemich, H. C. Rumsey, Jr. and L. R. Welch, "New upper bounds on the rate of a code via the Delsart-MacWilliams inequalities," *IEEE Trans. Inf. Theory*, vol 23, pp. 157-166 (1977).

[23] A. Ashikhmin, "Remarks on bounds for quantum codes," preprint quant-ph/9705037.

n	k	k'	d	d'	K	D
8	4	7	4	2	3	3
16	11	15	4	2	10	3
32	26	31	4	2	25	3
32	21	26	6	4	15	6
32	16	21	8	6	5	8
64	57	63	4	2	56	3
64	51	57	6	4	44	6
64	45	51	8	6	32	8
128	120	127	4	2	119	3
128	113	120	6	4	105	6
128	106	113	8	6	91	8
128	99	113	10	6	84	9
128	92	106	12	8	70	12
128	85	99	14	10	56	14
128	78	99	16	10	49	15
256	247	255	4	2	246	3
256	239	247	6	4	230	6
256	231	239	8	6	214	8
256	223	239	10	6	206	9
256	215	231	12	8	190	12
256	207	223	14	10	174	14
256	199	223	16	10	166	15

Table 1. Parameters $[[n, K, D]]$ of the quantum codes obtained from primitive binary BCH codes, for $n \leq 256$. The BCH codes have been extended by an overall parity check in order to allow the distance 3 quantum code to be obtained by combining a BCH code with the even weight code. For $D > 3$ if the unextended BCH codes are used, a $[[n-1, K+1, D-1]]$ quantum code is obtained.

n	k	k'	d	d'	K	D
22	15	21	4	2	14	3
22	12	15	6	4	5	6
46	33	45	4	2	32	3
46	29	33	6	4	16	6
52	43	51	4	2	42	3
74	64	73	4	2	63	3
74	55	64	6	4	45	4
74	46	55	10	6	27	9
86	77	85	4	2	76	3
86	69	77	6	4	60	6
90	78	89	4	2	77	3
90	67	78	6	4	55	6
90	56	67	10	6	33	9
90	45	56	12	10	11	12
94	83	93	4	2	82	3
94	78	83	6	4	67	6
94	68	78	8	6	52	8
94	58	78	10	6	42	9
94	53	68	12	8	27	12
106	93	104	4	2	92	3
106	81	93	6	4	68	6
106	75	81	8	6	50	8
106	71	81	10	6	46	9
118	105	117	4	2	104	3
118	93	105	6	4	80	6
118	81	93	8	6	56	8
118	69	93	10	6	44	9

Table 2. As table 1, but for non-primitive BCH codes with $n < 127$.